Effective Static Deadlock Detection

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Abstract

We present an effective static deadlock detection algorithm for Java. Our algorithm uses a novel combination of static analyses each of which approximates a different necessary condition for a deadlock. We have implemented the algorithm and report upon our experience applying it to a suite of multi-threaded Java programs. While neither sound nor complete, our approach is effective in practice, finding all known deadlocks as well as discovering previously unknown ones in our benchmarks with few false alarms.

1 Introduction

A deadlock in a shared-memory multi-threaded program is an unintended condition in which a set of threads blocks forever because each thread in the set is waiting to acquire a lock already held by another thread in the set. Today's concurrent programs are riddled with deadlocks—a problem further exacerbated by the shift to multicore processors. For instance, roughly 6500/198,000 ($\sim 3\%$) of the bug reports in Sun's bug database at `http://bugs.sun.com` involve the keyword "deadlock". Moreover, fixing other concurrency problems like races often involves introducing new synchronization, which in turn can introduce new deadlocks. Hence, *static deadlock detection* is valuable for testing and debugging such programs.

Previous static deadlock detection approaches are based on type systems [3,4], dataflow analyses [1,7,12,13,17,19, 21], or model checking [5,6,11] (Section 7). The annotation burden for type-based approaches is often significant while model checking approaches currently do not scale to beyond a few thousand lines of code. Approaches based on dataflow analysis, on the other hand, have been applied to large programs but are highly imprecise.

In this paper, we present an effective static deadlock detection algorithm for Java (Section 3). Conceptually, our algorithm considers every tuple $(t^a, l_1^a, l_2^a, t^b, l_1^b, l_2^b)$, where t^a, t^b denote abstract threads and l_1^a, l_2^a, l_1^b, l_2^b denote lock acquisition statements, and checks if any pair of threads abstracted by t^a and t^b may deadlock by waiting to acquire a pair of locks z_1 and z_2 at l_2^a and l_2^b, while already holding locks z_2 and z_1 at l_1^a and l_1^b. Our key idea is to express the complex property of deadlock freedom for a pair of threads/locks—a problem that no existing static analysis can directly solve effectively—in terms of six problems that can be solved effectively using existing static analyses:

- *reachable*: In some execution of the program, can a thread abstracted by t^a reach l_1^a and, after acquiring a lock at l_1^a, proceed to reach l_2^a while still holding the lock (and similarly for t^b, l_1^b, l_2^b)?
- *aliasing*: In some execution of the program, can a lock acquired at l_1^a be the same as a lock acquired at l_2^b (and similarly for l_2^a, l_1^b)?
- *escaping*: In some execution of the program, can a lock acquired at l_1^a be accessible from more than one thread (and similarly for each of l_2^a, l_1^b, and l_2^b)?
- *parallel*: In some execution of the program, can different threads abstracted by t^a and t^b simultaneously reach l_2^a and l_2^b, respectively?
- *non-reentrant*: In some execution of the program, can a thread abstracted by t^a acquire a lock at l_1^a it does not already hold and, while holding that lock, proceed to acquire a lock at l_2^a it does not already hold (and similarly for t^b, l_1^b, l_2^b)? If the thread acquires the same lock it already holds then the second lock acquisition cannot cause a deadlock as locks are *reentrant* in Java.
- *non-guarded*: In some execution of the program, can different threads abstracted by t^a and t^b reach l_1^a and l_1^b, respectively, without holding a *common* lock? If the two threads already hold a common lock then we call it a *guarding lock* (also called a *gate lock* [10]).

Each of these six necessary conditions is undecidable. Thus, any solution to each of them is necessarily unsound or incomplete. Our algorithm soundly approximates the first four conditions using well-known static analyses, namely, a *call-graph analysis*, a *may-alias analysis*, a *thread-escape* *analysis*, and a *may-happen-in-parallel analysis*, respectively. Soundly approximating the last two conditions, however, requires a *must-alias analysis*, which is much harder than may-alias analysis. We address this problem using a common unsound solution: we use our may-alias analysis to masquerade as a must-alias analysis—as a result, we may fail to report some real deadlocks.

We may also report false deadlocks, either due to imprecision in our approximation of the six conditions, or because the deadlock is prevented by some condition not considered by our algorithm (Section 6). However, our approach is extensible: additional conditions, perhaps specific to the language or even the application at hand, can easily be added. In fact, the *non-guarded* and *non-reentrant* conditions specifically target Java programs. These idioms, if not identified, cause any static deadlock detector for Java to report overwhelmingly many false deadlocks [10,21].

Our algorithm, while unsound and incomplete, is effective in practice. We have implemented it in a tool JADE (Section 4) and applied it to a suite of multi-threaded Java programs comprising over 1.5 MLOC. Our approach found all known deadlocks as well as discovered previously unknown ones in the suite, with few false alarms (Section 5).

2 Example

We first illustrate our approach on a real-world case: the JDK's logging facilities from package `java.util.logging`. These facilities are provided as a library whereas our approach uses whole-program static analyses and thus requires a *closed program*, i.e., a complete program with a main method. So the first step in applying our approach to an *open program* such as a library is to build a *harness* that simulates clients exercising the interface of the program. Currently, we construct harnesses manually. Our algorithm is not path-sensitive and it ignores the values of primitive data. Hence, it neither requires a detailed, fully concrete harness nor test input data.

A snippet of our harness for this example is shown in class `Harness` in Figure 1. For brevity, we omit access qualifiers on classes, methods, and fields. Also, we label object allocation sites h_1–h_3, synchronized methods m_1–m_3, and thread `run` methods m_4 and m_5. The harness creates and starts two threads which we identify by their object allocation sites h_1 and h_2. Thread h_1 calls static method `Logger.getLogger`, which returns the unique logger having the specified name, creating the logger if it does not already exist in the global logger manager. This manager, allocated at site h_3 and stored in static field `LogManager.manager`, maintains all existing loggers in a hashtable. On the other hand, thread h_2 calls instance method `LogManager.addLogger`, which adds the specified logger to the global logger manager's hashtable if it

```
class Harness {
     static void main(String[] args) {
11:     Thread v1 = new^h1 Thread() {
           void run^m4() {
13:           Logger.getLogger(...);
           }};
        v1.start();
16:     Thread v2 = new^h2 Thread() {
           void run^m5() {
18:           LogManager.manager.addLogger(...);
           }};
        v2.start();
     }
}

// snippet of java/util/logging/Logger.java
class Logger {
226: static sync^m1 Logger getLogger(String name) {
        LogManager lm = LogManager.manager;
228:    Logger l = lm.getLogger(name);
        if (l == null) {
           l = new Logger(...);
231:       lm.addLogger(l);
        }
        return l;
     }
}

// snippet of java/util/logging/LogManager.java
class LogManager {
     static final LogManager manager =
155:    new^h3 LogManager();
     Hashtable loggers = new Hashtable();
280: sync^m2 boolean addLogger(Logger l) {
        String name = l.getName();
        if (!loggers.put(name, l))
           return false;
        // ensure l's parents are instantiated
        for (...) {
           String pname = ...;
314:       Logger.getLogger(pname);
        }
        return true;
     }
420: sync^m3 Logger getLogger(String name) {
        return (Logger) loggers.get(name);
     }
}
```

Figure 1. Example Java program.

```
*** Stack trace of thread <Harness.java:11>:
LogManager.addLogger (LogManager.java:280)
  - this allocated at <LogManager.java:155>
  - waiting to lock {<LogManager.java:155>}
Logger.getLogger (Logger.java:231)
  - holds lock {<Logger.java:0>}
Harness$1.run (Harness.java:13)

*** Stack trace of thread <Harness.java:16>:
Logger.getLogger (Logger.java:226)
  - waiting to lock {<Logger.java:0>}
LogManager.addLogger (LogManager.java:314)
  - this allocated at <LogManager.java:155>
  - holds lock {<LogManager.java:155>}
Harness$2.run (Harness.java:18)
```

Figure 2. Example deadlock report.

does not already contain a logger with that name.

Our algorithm reports the *counterexample* shown in Figure 2 for this program. It is similar to a thread stack dump output by a dynamic tool except that it is produced by a static tool and hence may denote a false deadlock. To improve usability we provide additional details to help users determine whether the counterexample denotes a real deadlock or a false positive. First, although we cannot provide concrete addresses of threads, we can identify their allocation sites. For instance, the counterexample in Figure 2 reports a deadlock between threads h_1 and h_2, identified by allocation sites <Harness.java:11> and <Harness.java:16>, respectively. Likewise, instead of providing concrete addresses of locks, we provide abstract locks. An abstract lock is a set of *abstract objects* where each abstract object is a site at which the lock may be allocated. More generally, abstract objects may be sequences of multiple such sites, allowing different objects allocated at the same site to be distinguished (Section 3.1). In Figure 2, {<LogManager.java:155>} denotes the lock on the LogManager object allocated at site h_3 at LogManager.java:155 and stored in static field LogManager.manager while {<Logger.java:0>} denotes the lock on the implicitly allocated java.lang.Class object stored in the implicit static field Logger.class.

Finally, each instance method called in each stack trace is coupled with an abstract object denoting the site at which the distinguished this variable of that method is allocated. In Figure 2, instance method LogManager.addLogger in either stack trace is called in a context in which its this variable is allocated at site h_3 at LogManager.java:155. In more complex programs, the same method may be analyzed in multiple contexts (Section 3.1).

It is easy to see that the above counterexample denotes a real deadlock: thread h_1 waits to acquire the lock on LogManager.manager while holding the lock on Logger.class, whereas thread h_2 waits to acquire the lock on Logger.class while holding the lock on LogManager.manager.

Our algorithm reports another counterexample for the above program; the only difference is that the topmost call in the stack trace of thread h_1 is to method LogManager.getLogger from call site Logger.getLogger (Logger.java:228) instead of to method LogManager.addLogger from call site Logger.getLogger (Logger.java:231). Both methods attempt to acquire the same lock, on LogManager.manager, and hence both counterexamples denote the same deadlock. In our experience, our algorithm's ability to report all possible ways in which the same deadlock may occur helps in determining the best fix for the deadlock. In the above program, for instance, both counterexamples (and many similar ones resulting from parts of the program not shown here) contain the same last call in the stack trace of thread h_2, namely, code Logger.getLogger(pname) at LogManager.java:314. Indeed, the fix for this deadlock is to replace this code by the inlined body of method Logger.getLogger without its synchronization so that it does not hold the lock on Logger.class.

$$
\begin{array}{rrcl}
\textit{(method)} & m & \in & \mathbb{M} = \{\mathsf{m}_{main}, \mathsf{m}_{start}, ...\} \\
\textit{(local var.)} & v & \in & \mathbb{V} \\
\textit{(alloc. site)} & h & \in & \mathbb{H} \\
\textit{(list)} & [h_1 :: \ldots :: h_n] & \in & \mathbb{H}^n \\
\\
\textit{(abstract object)} & o & \in & \mathbb{O} = \mathbb{H}^0 \cup \mathbb{H}^1 \cup \mathbb{H}^2 \cup \ldots \\
\textit{(abstract context)} & c, t, l & \in & \mathbb{C} = \mathbb{O} \times \mathbb{M} \\
\textit{(synchronized argument)} & \mathsf{sync} & : & \mathbb{C} \rightharpoonup \mathbb{V} \\
\\
\textit{(call graph)} & \mathsf{cg} & \subseteq & (\mathbb{C} \times \mathbb{C}) \\
\textit{(points-to for locals)} & \mathsf{pt} & \subseteq & (\mathbb{C} \times \mathbb{V} \times \mathbb{O}) \\
\textit{(thread-escape)} & \mathsf{esc} & \subseteq & (\mathbb{C} \times \mathbb{V}) \\
\textit{(may-happen-in-parallel)} & \mathsf{mhp} & \subseteq & (\mathbb{C} \times \mathbb{C} \times \mathbb{C})
\end{array}
$$

Figure 3. Notation.

3 Algorithm

Our algorithm is based on sound and unsound approximations of our six necessary conditions (Section 3.2). Effectively approximating these conditions needs precise call-graph and points-to information—we use a form of combined call-graph and may-alias analysis called *k-object-sensitive analysis* [14] (Section 3.1). Finally, to improve usability, our algorithm generates and groups counterexamples to explain the deadlocks it detects (Section 3.3).

Before presenting our algorithm, we summarize our notation (Figure 3). Our algorithm takes as input a closed program with a main method denoted m_{main}. We use $\mathbb{M}$ to denote the set of all method implementations that may be reachable from m_{main}. $\mathbb{M}$ may be a crude over-approximation, e.g., one computed by Class Hierarchy Analysis (CHA). We use $\mathsf{m}_{start} \in \mathbb{M}$ to denote the start() method of class java.lang.Thread, the method used to explicitly spawn a thread. We use $\mathbb{V}$ to denote the set of all local variables referenced by methods in $\mathbb{M}$. We presume that each method $m \in \mathbb{M}$ may be synchronized on any one of its arguments, specified by $\mathsf{sync}((o, m))$ (o is irrelevant but simplifies our notation), but does not contain any other synchronized blocks in its body. If method m is not synchronized, then the partial function sync is not defined at (o, m). It is easy to transform any Java program to satisfy this restriction (Section 4).

Figure 3 also shows the relations produced by our four sound whole-program static analyses: call-graph analysis (cg), may-alias analysis (pt), thread-escape analysis (esc), and may-happen-in-parallel analysis (mhp). These relations are the ones we need to define our deadlock detector; internally our analyses track information in greater detail (e.g., the may-alias analysis tracks the contents of the heap and static fields). We outline how pt and cg are computed in Section 3.1; we reuse the thread-escape analysis and may-happen-in-parallel analysis from earlier work [15].

3.1 k-Object-Sensitive Analysis

k-object-sensitive analysis [14] is an *object sensitive*, *context sensitive*, and *flow insensitive* analysis that computes call-graph and points-to approximations.

The analysis is *object sensitive* in that it can represent different objects allocated at the same site by potentially different abstract objects. An abstract object $o \in \mathbb{O}$ is a finite sequence of object allocation sites denoted $[h_1 :: ... :: h_n]$. The first allocation site h_1 is the represented object's allocation site. The subsequent allocation sites $[h_2 :: ... :: h_n]$ represent the object denoted by the distinguished `this` variable of the method where o was allocated—thus, that `this` object was allocated at h_2, in a method whose `this` object is represented by $[h_3 :: ... :: h_n]$, and so on. For static methods, which lack the `this` variable, we represent the `this` object by $[]$ (which represents no objects).

The analysis is also *context sensitive* in that it can analyze each method implementation in potentially multiple abstract contexts. An abstract context $c \in \mathbb{C}$ is a pair (o, m) of an abstract object o and a method m such that o abstracts the `this` object of m; as above, for static methods $o = []$.

Finally, the analysis is *flow insensitive* as it computes global (instead of per program point) points-to information. This, however, does not adversely affect the precision of the analysis on local variables as our implementation operates on a Static Single Assignment (SSA) representation of the program (Section 4).

The analysis produces the following relations:

- cg $\subseteq (\mathbb{C} \times \mathbb{C})$, the context-sensitive call graph, contains each tuple $((o_1, m_1), (o_2, m_2))$ such that method m_1 may call method m_2 with its `this` object abstracted by o_2 when the `this` object of m_1 is abstracted by o_1.

- pt $\subseteq (\mathbb{C} \times \mathbb{V} \times \mathbb{O})$, the points-to information for local variables, contains each tuple (c, v, o) such that local variable v may point to abstract object o in abstract context c.

We illustrate how k-object-sensitive analysis computes these relations for our running example from Figure 1, concentrating on how object allocation sites and method call sites are handled. The analysis begins by deeming reachable the contexts of the main method $([], \texttt{Harness.main})$ and of every class initializer method (e.g., $([], \texttt{LogManager.<clinit>})$). As the analysis is flow insensitive, whenever a context (o, m) is reachable, every statement in the body of m is reachable.

The analysis presumes a positive integer associated with each object allocation site, called the k-value of that site. Consider any such site $v = \textbf{new}^h ...$, where $h \in \mathbb{H}$ and $v \in \mathbb{V}$, in a method m that the analysis deems reachable in a context (o, m). Then, the analysis adds tuple $((o, m), v, h \oplus_k o)$ to relation pt, where $h \oplus_k o$ is a finite non-empty sequence of object allocation sites whose head is h and whose tail comprises at most the $k - 1$ most significant sites in o in order. Our deadlock detection algorithm automatically chooses potentially different k-values for different sites (Section 4). For our running example, however, we presume $k = 1$ for all sites. The initially reachable context $([], \texttt{Harness.main})$ contains object allocation statements:

$$\texttt{v1} = \textbf{new}^{\mathsf{h_1}} \dots \text{ and } \texttt{v2} = \textbf{new}^{\mathsf{h_2}} \dots$$

and so the analysis adds the following tuples to pt:

$$(([], \texttt{Harness.main}), \texttt{v1}, [\mathsf{h_1}])$$
$$(([], \texttt{Harness.main}), \texttt{v2}, [\mathsf{h_2}])$$

If $n(\dots)$ is a static method call in a reachable context (o, m), the analysis adds tuple $((o, m), ([], n))$ to cg. Also, the analysis henceforth deems context $([], n)$ reachable.

If $v.n(\dots)$ is an instance method call, then the target method depends upon the run-time type of the object denoted by v. Every $((o, m), v, [h_1 :: ... :: h_n]) \in$ pt denotes a target in a potentially different context. The analysis thus adds $((o, m), ([h_1 :: ... :: h_n], n'))$ to cg, where n' is the target of a call to n for an object allocated at site h_1. We must also add $(([h_1 :: ... :: h_n], n'), \texttt{this}, [h_1 :: ... :: h_n])$ to pt—this treatment of the `this` argument is key to precision [14]. Also, the analysis henceforth deems context $([h_1 :: ... :: h_n], n')$ reachable. Furthermore, if $n' = \mathsf{m}_{start}$ (a thread is started), the context $([h_1 :: ... :: h_n], n'')$ is also deemed reachable, where n'' is the `run()` method of the class allocated at h_1.

For our running example, since the analysis has deemed context $([], \texttt{Harness.main})$ reachable and `Harness.main` contains calls `v1.start()` and `v2.start()`, the analysis adds the following tuples to cg:

$$(([], \texttt{Harness.main}), ([\mathsf{h_1}], \mathsf{m_{start}}))$$
$$(([], \texttt{Harness.main}), ([\mathsf{h_2}], \mathsf{m_{start}}))$$

and deems contexts $([\mathsf{h_1}], \mathsf{m_4})$ and $([\mathsf{h_2}], \mathsf{m_5})$ of the respective `run()` methods reachable.

3.2 Deadlock Computation

Our deadlock detection algorithm represents threads (t) by the abstract context of the thread's entry method, and

(Reachability) $c_1 \to c_2 \triangleright L$ iff $\exists n:\ c_1 \to^n c_2 \triangleright L$ where:
(1) $c \to^0 c \triangleright \emptyset$
(2) $c_1 \to^{n+1} c_2 \triangleright L'$ iff $\exists c, L:\ c_1 \to^n c \triangleright L \ \wedge\ (c, c_2) \in \mathsf{cg} \ \wedge\ L' = \begin{cases} L \cup \{c\} & \text{if } \mathsf{sync}(c) \text{ defined} \\ L & \text{otherwise} \end{cases}$

(Lock Aliasing) $\mathsf{mayAlias}(l_1, l_2)$ iff $\exists o:\ (l_1, \mathsf{sync}(l_1), o) \in \mathsf{pt} \ \wedge\ (l_2, \mathsf{sync}(l_2), o) \in \mathsf{pt}$
(Lock-set Aliasing) $\mathsf{mayAlias}(L_1, L_2)$ iff $\exists l_1 \in L_1, l_2 \in L_2:\ \mathsf{mayAlias}(l_1, l_2)$

Figure 4. Reachability, locks, and aliasing.

lock acquisitions (l) by the abstract context of synchronized methods; the latter suffices as we presume that methods do not contain any synchronized blocks in their body. We represent sets of held locks (L) by sets of abstract contexts of synchronized methods that acquire the corresponding locks.

Figure 4 defines some properties of threads, lock acquisitions, and lock sets that we derive from pt and cg and use in the rest of our algorithm. We use $c_1 \to c_2 \triangleright L$ to denote that context c_2 may be reachable from context c_1 along some path in some thread, and, moreover, a thread executing that path may hold set of locks L upon reaching c_2 (we elide $\triangleright L$ when the locks are irrelevant). We use $\mathsf{mayAlias}(l_1, l_2)$ to denote that lock acquisitions at l_1 and l_2 may acquire the same lock. We extend mayAlias to lock sets as usual.

We use reachability ($\to$) to approximate the set of startable threads and reachable lock acquisitions:

$$
\begin{aligned}
\mathsf{threads} &= \{\, x \mid \exists n:\ x \in \mathsf{threads}_n \,\} \text{ where} \\
\mathsf{threads}_0 &= \{\, ([], \mathsf{m}_{main}) \,\} \\
\mathsf{threads}_{n+1} &= \mathsf{threads}_n \cup \\
&\quad \{\, (o, \mathtt{run}) \mid c \in \mathsf{threads}_n \wedge c \to (o, \mathsf{m}_{start}) \,\} \\
\mathsf{locks} &= \{\, c \mid c' \in \mathsf{threads} \wedge c' \to c \wedge \mathsf{sync}(c) \text{ defined} \,\}
\end{aligned}
$$

For our running example, we have:

$$
\begin{gathered}
\mathsf{threads} = \{\, ([], \mathsf{m}_{main}), \mathsf{t}_1, \mathsf{t}_2 \,\} \\
\mathsf{locks} = \{\, \mathsf{l}_1, \mathsf{l}_2, \mathsf{l}_3 \,\} \\
\mathsf{t}_1 \triangleq ([\mathsf{h}_1], \mathsf{m}_4) \quad \mathsf{t}_2 \triangleq ([\mathsf{h}_2], \mathsf{m}_5) \\
\mathsf{l}_1 \triangleq ([], \mathsf{m}_1) \quad \mathsf{l}_2 \triangleq ([\mathsf{h}_3], \mathsf{m}_2) \quad \mathsf{l}_3 \triangleq ([\mathsf{h}_3], \mathsf{m}_3)
\end{gathered}
$$

A deadlock six-tuple $d = (t^a, l_1^a, l_2^a, t^b, l_1^b, l_2^b)$ denotes a deadlock involving a pair of locks z_1 and z_2 such that thread t^a holds lock z_1 it acquired at synchronized method l_1^a and is waiting to acquire lock z_2 at l_2^a while thread t^b holds lock z_2 it acquired at synchronized method l_1^b and is waiting to acquire lock z_1 at l_2^b. Conceptually, our deadlock detection algorithm simply filters all potential deadlocks through our six necessary conditions, computing the final set of potential deadlocks to be reported as:

$$
\begin{aligned}
\mathsf{finalDeadlocks} = \{\, & d \mid d = (t^a, l_1^a, l_2^a, t^b, l_1^b, l_2^b) \wedge \\
& t^a, t^b \in \mathsf{threads} \ \wedge\ l_1^a, l_2^a, l_1^b, l_2^b \in \mathsf{locks} \wedge \\
& \mathsf{reachableDeadlock}\ d \ \wedge\ \mathsf{aliasingDeadlock}\ d \wedge \\
& \mathsf{escapingDeadlock}\ d \ \wedge\ \mathsf{parallelDeadlock}\ d \wedge \\
& \mathsf{nonReentDeadlock}\ d \ \wedge\ \mathsf{nonGrdedDeadlock}\ d \,\}
\end{aligned}
$$

Our six necessary conditions are formally defined in Sections 3.2.1–3.2.6 below.

For our running example, we focus on two potential deadlocks:

$$
\begin{aligned}
\mathsf{d}_1 &\triangleq (\mathsf{t}_1, \mathsf{l}_1, \mathsf{l}_2, \mathsf{t}_2, \mathsf{l}_2, \mathsf{l}_1) \\
\mathsf{d}_2 &\triangleq (\mathsf{t}_1, \mathsf{l}_1, \mathsf{l}_3, \mathsf{t}_2, \mathsf{l}_2, \mathsf{l}_1)
\end{aligned}
$$

Each of these tuples denotes a possible deadlock between abstract threads t_1 and t_2. In both tuples, thread t_2 holds a lock at l_2 (context $([\mathsf{h}_3], \mathsf{m}_2)$) and is waiting to acquire a lock at l_1 (context $([], \mathsf{m}_1)$). Also, in both tuples, thread t_1 holds a lock at l_1, but it is waiting to acquire a lock at l_2 in tuple d_1 and a lock at l_3 (context $([\mathsf{h}_3], \mathsf{m}_3)$) in tuple d_2. As we will see, d_1 and d_2 pass all six conditions and are thus contained in finalDeadlocks.

3.2.1 Computation of reachableDeadlock

For a tuple $(t^a, l_1^a, l_2^a, t^b, l_1^b, l_2^b)$ to be a deadlock our *reachable* condition must be satisfied: Can a thread abstracted by t^a reach l_1^a and, after acquiring a lock at l_1^a, proceed to reach l_2^a while still holding the lock (and similarly for t^b, l_1^b, l_2^b)?

Our algorithm uses the reachability property (Figure 4) to approximate this condition:

$$
\begin{gathered}
\mathsf{reachableDeadlock}\ (t^a, l_1^a, l_2^a, t^b, l_1^b, l_2^b) \ \text{ if} \\
t^a \to l_1^a \ \wedge\ l_1^a \to l_2^a \ \wedge\ t^b \to l_1^b \ \wedge\ l_1^b \to l_2^b
\end{gathered}
$$

For our running example, it is easy to see that thread t_1 reaches l_1, then l_3 and subsequently l_2, while t_2 reaches l_2 and then l_1. Thus both tuples d_1 and d_2 satisfy reachableDeadlock.

3.2.2 Computation of aliasingDeadlock

For a tuple $(t^a, l_1^a, l_2^a, t^b, l_1^b, l_2^b)$ to be a deadlock our *aliasing* condition must be satisfied: Can a lock acquired at l_1^a be the same as a lock acquired at l_2^b (and, similarly for l_2^a, l_1^b)? Our algorithm uses the mayAlias property (Figure 4) to approximate this condition:

$$
\begin{gathered}
\mathsf{aliasingDeadlock}\ (t^a, l_1^a, l_2^a, t^b, l_1^b, l_2^b) \ \text{ if} \\
\mathsf{mayAlias}(l_1^a, l_2^b) \ \wedge\ \mathsf{mayAlias}(l_2^a, l_1^b)
\end{gathered}
$$

For our running example, both tuples d_1 and d_2 satisfy aliasingDeadlock: predicates $\mathsf{mayAlias}(\mathsf{l}_1, \mathsf{l}_1)$ and

$\mathsf{mayAlias}(\mathsf{l}_2, \mathsf{l}_2)$ hold trivially and hence tuple d_1 satisfies aliasingDeadlock; additionally, $\mathsf{mayAlias}(\mathsf{l}_3, \mathsf{l}_2)$ holds because abstract object $[\mathsf{h}_3]$ satisfies the two conjuncts in the definition of mayAlias, and hence tuple d_2 also satisfies aliasingDeadlock.

3.2.3 Computation of escapingDeadlock

The JDK contains many classes (e.g. `java.util.Vector`) with synchronized methods. When such objects cannot be accessed by more than one thread, they cannot participate in a deadlock. Thus, for a tuple (t^a, l_1^a, l_2^a, t^b, l_1^b, l_2^b) to be a deadlock our *escaping* condition must be satisfied: Can a lock acquired at l_1^a be accessible from more than one thread (and similarly for each of l_2^a, l_1^b, l_2^b)?

We approximate this condition using a thread-escape analysis. Our application of this analysis to static deadlock detection appears novel and we quantify the need for it in our experiments (Section 5).

The thread-escape problem is usually defined as follows: "In some execution, is some object allocated at a given site h accessible from more than one thread?" To increase precision, we refine the notion of thread-escape to track *when* an object escapes. This allows the *escaping* condition to eliminate some deadlock reports on objects that later escape to other threads. Formally, (c, v) must be in relation esc if argument v of abstract context c may be accessible from more than one thread. Our *escaping* condition is thus:

$$\begin{array}{l}\mathsf{escapingDeadlock}\ (t^a, l_1^a, l_2^a, t^b, l_1^b, l_2^b)\ \text{ if} \\ \quad (l_1^a, \mathsf{sync}(l_1^a)) \in \mathsf{esc}\ \wedge\ (l_2^a, \mathsf{sync}(l_2^a)) \in \mathsf{esc}\ \wedge \\ \quad (l_1^b, \mathsf{sync}(l_1^b)) \in \mathsf{esc}\ \wedge\ (l_2^b, \mathsf{sync}(l_2^b)) \in \mathsf{esc}\end{array}$$

For our running example, `LogManager.manager` (l_2, l_3) and `Logger.class` (l_1), being static fields, clearly escape everywhere, and so both tuples d_1 and d_2 satisfy escapingDeadlock.

3.2.4 Computation of parallelDeadlock

For a tuple (t^a, l_1^a, l_2^a, t^b, l_1^b, l_2^b) to be a deadlock our *parallel* condition must be satisfied: Can *different* threads abstracted by t^a and t^b *simultaneously* reach l_2^a and l_2^b, respectively? The motivation for checking this condition is two-fold. First, it eliminates each tuple $(t, *, *, t, *, *)$ where t abstracts at most one thread in any execution. The most common example of such an abstract thread is $([], \mathsf{m}_{main})$, but it also applies to any thread class allocated at most once in every execution. The second motivation is that even if different threads abstracted by t^a and t^b may be able to reach l_2^a and l_2^b, respectively, the thread structure of the program may forbid them from doing so *simultaneously*, namely, threads t^a and t^b may be in a "parent-child" relation, causing l_2^a to *happen before* l_2^b in all executions.

We approximate these two conditions using a may-happen-in-parallel analysis that computes relation mhp which contains each tuple $(t_1, (o, m), t_2)$ such that a thread abstracted by t_2 may be running in parallel when a thread abstracted by t_1 reaches method m in context o. Our may-happen-in-parallel analysis is simple and only models the program's thread structure, ignoring locks and other kinds of synchronization (fork-join, barrier, etc). Our *parallel* condition is thus:

$$\begin{array}{l}\mathsf{parallelDeadlock}\ (t^a, l_1^a, l_2^a, t^b, l_1^b, l_2^b)\ \text{ if} \\ \quad (t^a, l_2^a, t^b) \in \mathsf{mhp}\ \wedge\ (t^b, l_2^b, t^a) \in \mathsf{mhp}\end{array}$$

For our running example, clearly nothing prevents t_1 and t_2 from running in parallel, so tuples d_1 and d_2 satisfy parallelDeadlock.

3.2.5 Computation of nonReentDeadlock

In Java, a thread can re-acquire a lock it already holds. This *reentrant* lock acquisition cannot cause a deadlock. Thus, for a tuple (t^a, l_1^a, l_2^a, t^b, l_1^b, l_2^b) to be a deadlock our *non-reentrant* condition must be satisfied: Can a thread abstracted by t^a acquire a lock at l_1^a it does not already hold and, while holding that lock, proceed to acquire a lock at l_2^a it does not already hold (and similarly for t^b, l_1^b, l_2^b)?

Soundly identifying reentrant locks requires must-alias analysis. Must-alias analysis, however, is much harder than may-alias analysis. Instead, we use our may-alias analysis itself to unsoundly check that whenever a thread abstracted by t acquires a lock at l_1 and, while holding that lock, proceeds to acquire a lock at l_2, then the lock it acquires at l_1 or l_2 *may* (soundness requires *must*) be already held by the thread—a property approximated by reentrant:

$$\begin{array}{l}\mathsf{reentrant}(t, l_1, l_2)\ \text{ iff }\ l_1 = l_2 \vee \\ \quad (\forall L_1 :\ (t \rightarrow l_1 \triangleright L_1 \implies \mathsf{mayAlias}(\{l_1, l_2\}, L_1))) \vee \\ \quad (\forall L_2 :\ (l_1 \rightarrow l_2 \triangleright L_2 \implies \mathsf{mayAlias}(\{l_2\}, L_2)))\end{array}$$

Intuitively, the first conjunct checks that the locks acquired at l_1 and l_2 may be the same. The second conjunct checks that when a thread abstracted by t reaches up to but not including l_1, the set of locks L_1 it holds may contain the lock it will acquire at l_1 or l_2. The third conjunct checks that when the thread proceeds from l_1 and reaches up to but not including l_2, the set of locks L_2 it holds may contain the lock it will acquire at l_2. Next, we use the reentrant predicate to approximate our *non-reentrant* condition as follows:

$$\begin{array}{l}\mathsf{nonReentDeadlock}\ (t^a, l_1^a, l_2^a, t^b, l_1^b, l_2^b)\ \text{ if} \\ \quad \neg\mathsf{reentrant}(t^a, l_1^a, l_2^a)\ \wedge\ \neg\mathsf{reentrant}(t^b, l_1^b, l_2^b)\end{array}$$

The above approximation itself is sound but the approximation performed by the reentrant predicate it uses is unsound; thus, a tuple that does not satisfy nonReentDeadlock is not provably deadlock-free.

For our running example, the two locks acquired by either thread do not alias, and no locks are acquired prior to the first lock or between the first and second lock in either thread, so tuples d_1 and d_2 satisfy nonReentDeadlock.

3.2.6 Computation of nonGrdedDeadlock

One approach to preventing deadlock is to acquire a common *guarding lock* in all threads that might deadlock. Thus, for a tuple (t^a, l^a_1, l^a_2, t^b, l^b_1, l^b_2) to be a deadlock our *non-guarded* condition must be satisfied: Can threads abstracted by t^a and t^b reach l^a_1 and l^b_1, respectively, without already holding a *common* lock?

Soundly identifying guarding locks, like reentrant locks, needs a must-alias analysis. We once again use our may-alias analysis to unsoundly check whether every pair of threads abstracted by t^a and t^b *may* (soundness requires *must*) hold a common lock whenever they reach l^a and l^b, respectively—a property approximated by guarded:

$$\begin{array}{l}\mathsf{guarded}(t^a, l^a, t^b, l^b) \text{ iff } \forall L_1, L_2 : \\ \quad (t^a \to l^a \triangleright L_1 \ \wedge\ t^b \to l^b \triangleright L_2) \implies \mathsf{mayAlias}(L_1, L_2)\end{array}$$

Then, we use the guarded predicate to approximate our *non-guarded* condition as follows:

$$\begin{array}{l}\mathsf{nonGrdedDeadlock}\ (t^a, l^a_1, l^a_2, t^b, l^b_1, l^b_2) \text{ if} \\ \quad \neg\, \mathsf{guarded}(t^a, l^a_1, t^b, l^b_1)\end{array}$$

The above approximation itself is sound but the approximation performed by the guarded predicate it uses is unsound; thus, a tuple that does not satisfy nonGrdedDeadlock is not necessarily deadlock-free.

For our running example, as we saw for nonReentDeadlock, no locks are acquired prior to the first lock, so tuples d_1 and d_2 satisfy nonGrdedDeadlock.

3.3 Post-Processing

Our algorithm reports a *counterexample* for each tuple in finalDeadlocks. The counterexample reported for a tuple (t^a, l^a_1, l^a_2, t^b, l^b_1, l^b_2) in finalDeadlocks consists of a pair of paths P_1 and P_2 in the context-sensitive call graph denoting possible stack traces of threads abstracted by t^a and t^b, respectively, at the point of the deadlock. Specifically, P_1 is the shortest path from t^a to l^a_2 via l^a_1 and, similarly, P_2 is the shortest path from t^b to l^b_2 via l^b_1. Unlike stack traces reported by a dynamic tool, however, paths P_1 and P_2 may be infeasible; we aid the user in comprehending them by providing additional details such as the context in which each instance method along each of those paths is called and the set of abstract objects of each lock that is synchronized along each of those paths. For our running example, Figure 2 shows the counterexample reported for tuple d_1.

Our algorithm also groups together counterexamples likely to be symptoms of the same deadlock. For each tuple (t^a, l^a_1, l^a_2, t^b, l^b_1, l^b_2) in finalDeadlocks, it computes a pair of lock types (τ_1, τ_2) as the least upper bounds of the types of abstract objects in the points-to sets O_1 and O_2 of the two locks involved in the deadlock where:

$$\begin{array}{lcl}O_1 &=& \{\ o \mid (l^a_1, \mathsf{sync}(l^a_1), o) \in \mathsf{pt} \ \wedge\ (l^b_2, \mathsf{sync}(l^b_2), o) \in \mathsf{pt}\ \} \\ O_2 &=& \{\ o \mid (l^a_2, \mathsf{sync}(l^a_2), o) \in \mathsf{pt} \ \wedge\ (l^b_1, \mathsf{sync}(l^b_1), o) \in \mathsf{pt}\ \}\end{array}$$

Then, our algorithm groups together the counterexamples reported for tuples in finalDeadlocks that have the same pair of lock types. For our running example, both tuples d_1 and d_2 have the same pair of lock types (`java.lang.Class`, `java.util.logging.LogManager`). Hence, our algorithm groups their counterexamples together.

4 Implementation

We implemented our deadlock detection algorithm in a tool called JADE. JADE takes as input a closed Java program in bytecode form and, optionally, as source code (the latter is used only to report source-level counterexamples). It uses the Soot framework [18] to construct a 0-CFA-based call graph to determine the set $\mathbb{M}$ of all methods that may be reachable from the main method. It rewrites each synchronized block `synchronized (v) { s }` as a call to a fresh static method, synchronized on argument `v` with body `s`. It then converts the program into Static Single Assignment (SSA) form to increase the precision of the flow-insensitive k-object-sensitive analysis.

JADE then uses the results of k-object-sensitive analysis to perform the thread-escape and may-happen-in-parallel analyses. All three analyses are expressed in Datalog and solved using `bddbddb` [20], a Binary Decision Diagram (BDD)-based Datalog solver. BDDs compactly represent the input relations, such as those representing basic facts about the program (e.g., function sync), as well as the relations output by these analyses (e.g., pt, mhp, etc.). Finally, JADE runs our deadlock detection algorithm which is also expressed in Datalog and computes relation finalDeadlocks that approximates the set of tuples satisfying our six necessary conditions for a deadlock.

Our implementation of k-object-sensitive analysis is parameterized by three parameters:

- $\mathcal{M} \subseteq \mathbb{M}$ containing each method that must be analyzed context-insensitively (i.e., in the lone context []).
- $\mathcal{V} \subseteq \mathbb{V}$ containing each local variable whose points-to information must be maintained context-insensitively (i.e., in the lone context []).
- $\mathcal{K} : \mathbb{H} \to \mathbb{N}^+$ mapping each object allocation site to a positive integer called its k-value (Section 3.1).

For scalability, our k-object-sensitive analysis uses an iterative refinement-based approach: we run all three analyses and the deadlock detection algorithm in each iteration using increasingly refined $\mathcal{M}$, $\mathcal{V}$, and $\mathcal{K}$. In the first iteration, the cheapest possible k-object-sensitive analysis is run, using $\mathcal{M} = \mathbb{M}$, $\mathcal{V} = \mathbb{V}$ and $\mathcal{K} = \lambda h.1$, which is effectively a 0-CFA-based analysis, and finalDeadlocks is computed. These deadlocks, however, typically contain many false positives due to the imprecision of 0-CFA-based analysis (Section 5). Hence, instead of being reported, they are used to refine parameters $\mathcal{M}$, $\mathcal{V}$, and $\mathcal{K}$ and the k-object-sensitive analysis is re-run. The refinement algorithm considers each tuple in finalDeadlocks as an *effect* of imprecision and finds all its possible *causes* in terms of $\mathcal{M}$, $\mathcal{V}$, and $\mathcal{K}$ (e.g., a certain method must be analyzed context-sensitively, the k-value of a certain site must be increased, etc.). Since the other analyses depend upon k-object-sensitive analysis, they are also re-run, and finally the deadlock detection algorithm itself is re-run to compute a new finalDeadlocks. The process terminates either if finalDeadlocks is empty or if its size begins to grow; the latter termination criterion prevents overwhelming the user with too many reports denoting the same deadlock.

5 Experiments

We evaluated JADE on a suite of multi-threaded Java programs comprising over 1.5 MLOC. The suite includes the multi-threaded benchmarks from the Java Grande suite (`moldyn`, `montecarlo`, and `raytracer`); from ETH, a Traveling Salesman Problem implementation (`tsp`), a successive over-relaxation benchmark (`sor`) and a web crawler (`hedc`); a website download and mirror tool (`weblech`); a web spider engine (`jspider`); W3C's web server platform (`jigsaw`); and Apache's FTP server (`ftp`). The suite also includes open programs for which we manually wrote harnesses: Apache's database connection pooling library (`dbcp`); a fast caching library (`cache4j`); the JDK4 logging facilities (`logging`); and JDK4 implementations of lists, sets, and maps wrapped in synchronized collections (`collections`). All our benchmarks along with JADE's deadlocks reports are available at `http://chord.stanford.edu/deadlocks.html`

Table 1 summarizes JADE's results. The 'LOC', 'classes', 'methods', and 'syncs' columns show the numbers of lines of code, classes, methods, and synchronized statements deemed reachable from the main method by Soot's 0-CFA-based analysis. The 'time' column provides the total running time of JADE. The experiments were performed on a 64-bit Linux server with two 2GHz Intel Xeon quad-core processors and 8GB memory. JADE, however, is single-threaded and 32-bit, and hence utilizes only a single core and at most 4GB memory.

The '0-CFA' and 'k-obj.' columns give the size of finalDeadlocks after one and two iterations of our algorithm (Section 4)—finalDeadlocks is empty or starts to grow, and JADE terminates, after at most two iterations for all our benchmarks. The first iteration uses a k-object-sensitive analysis that is essentially a 0-CFA-based analysis (Section 4). The difference between the two columns, most notable for `hedc`, `weblech`, `jspider`, `ftp`, and `dbcp`, is the number of extra false positives that would be reported by a 0-CFA-based analysis over a k-object-sensitive one. All previous static deadlock detectors we are aware of employ a 0-CFA-based analysis or an even more imprecise CHA-based analysis; moreover, they exclude checking one or more of our six necessary conditions (Section 7).

Figure 5 justifies the need for the *escaping*, *parallel*, *non-reentrant* and *non-guarded* conditions—we consider the *reachable* and *aliasing* conditions fundamental to our deadlock definition. We measure the effectiveness of a particular condition by switching it off and observing the increase in the size of finalDeadlocks. The graphs exclude benchmarks `moldyn`, `raytracer`, `sor`, and `cache4j` as the size of finalDeadlocks is not noticeably affected for any of them by switching off any single condition—these benchmarks are relatively small and have a relatively simple synchronization structure (indicated by the numbers in the 'sync' column in Table 1) and gain no significant benefit from any one condition.

The left graph in Figure 5 shows the effectiveness of the sound *escaping* and *parallel* conditions. The bars are normalized to the number of deadlocks obtained by checking only the *reachable* and *aliasing* conditions. The 'soundDeadlocks' partition of each bar denotes the number of deadlocks obtained by checking all four sound conditions. The 'only Par.' (resp. 'only Esc.') partition denotes the number of deadlocks that are soundly filtered out exclusively by *parallel* (resp. *escaping*). The 'Esc. or Par.' partition denotes the number of deadlocks that are filtered out by both *parallel* and *escaping*. The right graph in Figure 5 shows the effectiveness of the unsound *non-reentrant* and *non-guarded* conditions. The bar for each benchmark in this graph further partitions the 'soundDeadlocks' partition of the bar for the corresponding benchmark in the left graph. The 'finalDeadlocks' partition denotes the size of finalDeadlocks. The 'only N.G.' (resp. 'only N.R.') partition denotes the number of deadlocks that are filtered out exclusively by *non-guarded* (resp. *non-reentrant*). Finally, the 'N.R. or N.G.' partition denotes the number of deadlocks that are filtered out by both *non-guarded* and *non-reentrant*. In summary, we see that each condition is important for some benchmark.

Our algorithm generates a counterexample for each tuple in finalDeadlocks reported under column 'k-obj.' in Table 1. These counterexamples are grouped by the pair of

benchmark	benchmark size				time	finalDeadlocks		lock type pairs	
	LOC	classes	methods	syncs		0-CFA	k-obj.	total	real
`moldyn`	31,917	63	238	12	4m48s	0	0	0	0
`montecarlo`	157,098	509	3447	190	7m53s	0	0	0	0
`raytracer`	32,576	73	287	16	4m51s	0	0	0	0
`tsp`	154,288	495	3335	189	7m48s	0	0	0	0
`sor`	32,247	57	208	5	4m48s	0	0	0	0
`hedc`	160,071	530	3552	204	21m15s	7,552	2,358	22	19
`weblech`	184,098	656	4620	238	32m09s	4,969	794	22	19
`jspider`	159,494	557	3595	205	15m34s	725	4	1	0
`jigsaw`	154,584	497	3346	184	15m23s	23	18	3	3
`ftp`	180,904	642	4383	252	35m55s	16,259	3,020	33	24
`dbcp`	168,018	536	3602	227	16m04s	320	16	4	3
`cache4j`	34,603	72	218	7	4m43s	0	0	0	0
`logging`	167,923	563	3852	258	9m01s	4,134	4,134	98	94
`collections`	38,961	124	712	55	5m42s	598	598	16	16

Table 1. Experimental results.

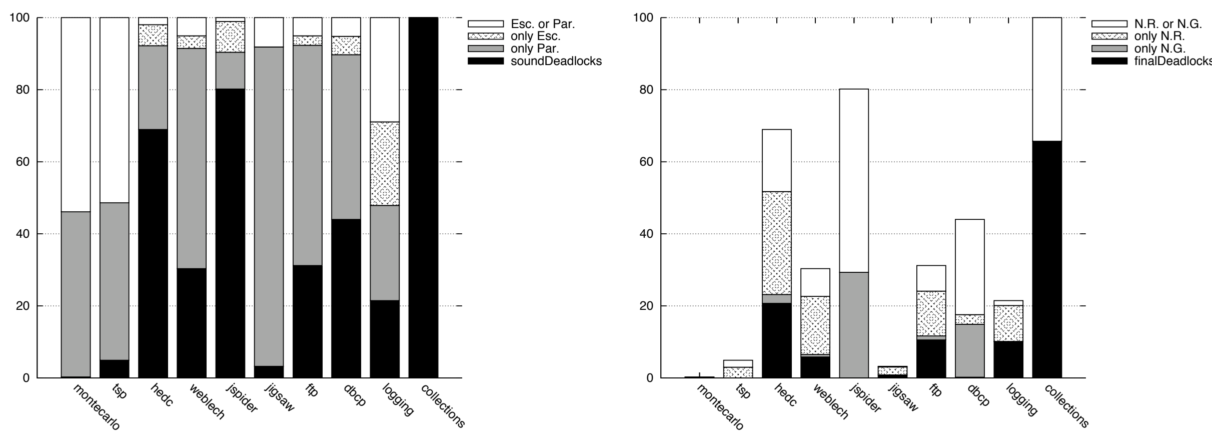

Figure 5. Contributions of individual analyses.

types of the locks involved in the deadlock (Section 3.3). The 'total' column in the table denotes the total number of such groups for each benchmark. The last 'real' column denotes the number of groups which contain at least one real deadlock. We confirmed real deadlocks by creating concrete test cases that were able to exhibit them. The deadlocks in `hedc`, `weblech`, `jigsaw`, and `ftp` were not in application code but in the JDK's logging facilities implemented in `java.util.logging`. This was our primary motivation for studying the `logging` benchmark; all the deadlocks reported for the above benchmarks are also included in those reported for `logging`. Additionally, this benchmark includes the previously known deadlock that is explained in our running example (Section 2) but is not triggered by any of the other benchmarks.

We found three application-level deadlocks in `dbcp` of which one was previously known. Finally, all deadlocks reported in benchmark `collections` are real and previously known. Strictly speaking, these are not bugs in the JDK collections per se but they indicate ways in which clients could erroneously use those collections and trigger deadlocks. We included `collections`, studied in previous work on deadlock detection [21], to confirm that our unsound approach could find all known deadlocks.

6 Limitations

Our deadlock detection algorithm is unsound. We begin by noting that it only reports deadlocks between two threads and two locks. Deadlocks between more than two threads/locks are possible and it is easy at least in principle to extend our approach to detect such deadlocks. However, empirical evidence from bug databases of popular open-source Java programs, such as `http://bugs.sun.com` and `http://issues.apache.org`, shows that the vast majority of deadlocks involve only two threads/locks (in fact, we did not encounter a single deadlock involving more than two threads/locks in perusing the above databases).

Our algorithm detects reentrant locks and guarding locks unsoundly (Sections 3.2.5 and 3.2.6). Two promising future directions are to check our *non-reentrant* condition using the form of must-alias analysis used to check finite-state properties [8] and to check our *non-guarded* condition using the form of must-alias analysis used to check races [16].

The key source of false positives in our experiments is the relatively imprecise thread-escape analysis used by our algorithm. Existing work on this analysis was driven primarily by the need to eliminate redundant synchronization in Java programs and subsided in recent years after modern JVMs diminished the run-time speedups achieved by this optimization. We hope our application of this analysis to static deadlock detection, and in our earlier work to static race detection [15], will renew advances in this analysis.

Our algorithm only detects deadlocks due to lock-based synchronization whereas other kinds of synchronization, notably wait-notify in Java, can cause deadlocks as well which our algorithm does not report.

Finally, our implementation ignores the effects of native methods and reflection in Java though we mitigate this problem by manually providing "stubs" for common native methods and annotations for statically resolving dynamic class loading sites in the JDK library.

7 Related work

Previous work on deadlock detection for shared-memory multi-threaded programs includes static approaches based on type systems, dataflow analysis, or model checking, as well as dynamic approaches.

7.1 Type Systems

Boyapati et al. [3, 4] present an ownership type system for Java that allows programmers to specify a partial order among locks. The type checker statically ensures that well-typed programs are deadlock-free. Our approach is unsound and cannot prove deadlock freedom. On the other hand, it does not require annotations and scales to larger programs.

7.2 Dataflow Analysis

Artho and Biere [1] augment Jlint, a static dataflow analysis based bug-finding tool for Java, with checks for several patterns that could indicate deadlocks. It performs local (per class or per method) analyses and cannot, for instance, infer that syntactically different expressions or synchronized blocks in methods of different classes may hold the same lock. Jlint lies in the category of lightweight tools that are unsound and incomplete but target common bug patterns and scale well; another similar tool is LockLint for C [17].

Von Praun [19] presents an algorithm that performs whole-program 0-CFA-based call-graph and may-alias analysis of Java programs to compute the static lock-order graph and reports cycles in it as possible deadlocks. Unlike our approach, his algorithm can report deadlocks involving more than two threads/locks. Like our approach, however, his algorithm is unsound and incomplete, and it checks necessary conditions for a deadlock that amount to our *reachable*, *aliasing*, and *non-reentrant* conditions, but not our *parallel*, *escaping*, and *non-guarded* conditions.

Williams et al. [21] present an algorithm that traverses the given Java program's call graph bottom-up and builds a lock-order graph summary per method. It then merges the summaries of thread entry methods into a global lock-order graph by unifying may-aliasing lock nodes together, and reports cycles in it as potential deadlocks. Unlike our approach, their algorithm can report deadlocks involving more than two threads/locks. Also, unlike our unsound checking of the *non-reentrant* condition, they handle reentrant locks soundly, but only detect them when lock expressions are local variables (as opposed to fields). This coupled with their CHA-based call-graph and may-alias analysis (which is less precise than a 0-CFA-based one) and the lack of checking of the *parallel*, *escaping*, and *non-guarded* conditions leads to significant imprecision which they address by applying several unsound heuristics.

Engler and Ashcraft [7] present RacerX, a static tool that performs flow-sensitive interprocedural analysis of C programs to compute the static lock-order graph and reports cycles in it as possible deadlocks. Their approach scales well but is highly imprecise and employs heuristics for ranking the deadlock reports in decreasing order of likelihood.

Masticola et al. [12, 13] present sound deadlock detection algorithms for various parallelism and synchronization models mainly in the context of Ada. A key aspect of their approach is *non-concurrency analysis* which may be viewed as the counterpart of our may-happen-in-parallel analysis.

7.3 Model Checking

The SPIN model checker has been used to verify deadlock freedom for Java programs by translating them into Promela, SPIN's modeling language [6, 11]. Model checking based on counterexample-guided abstraction refinement has also been applied to deadlock detection in message passing based C programs [5]. A general limitation of model checking approaches is that they presume that the input program has a finite and tractable state-space.

7.4 Dynamic Analysis

While deadlocks actually occurring in executions are easy to detect, dynamic approaches such as Visual Threads [9] monitor the order in which locks are held by each thread in an execution and report cycles in the resulting dynamic lock-order graph as potential deadlocks that could occur in a different execution. The Goodlock algorithm [2,10] extends this approach to reduce false positives, namely, it tracks thread fork/join events and guarding locks that render cycles infeasible; this is akin to checking our *parallel* and *non-guarded* conditions, respectively. Like any dynamic analysis, these approaches are inherently unsound and cannot be applied to open programs and without test input data.

8 Conclusion

We have presented a novel static deadlock detection algorithm for Java that uses four static analyses to approximate six necessary conditions for a deadlock. We have implemented and applied it to a suite of multi-threaded Java programs comprising over 1.5 MLOC. While unsound and incomplete, our approach is effective in practice, finding all known deadlocks as well as discovering previously unknown ones in our benchmarks with few false alarms.

9 Acknowledgment

We would like to thank Pallavi Joshi, Christos Stergiou, and the anonymous reviewers for their valuable comments. This research was supported in part by a generous gift from Intel, by Microsoft and Intel funding (award #20080469) and by matching funding by U.C. Discovery (award #DIG07-10227).

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