



SEJITS: Raising the Abstraction Level of Productivity Programming

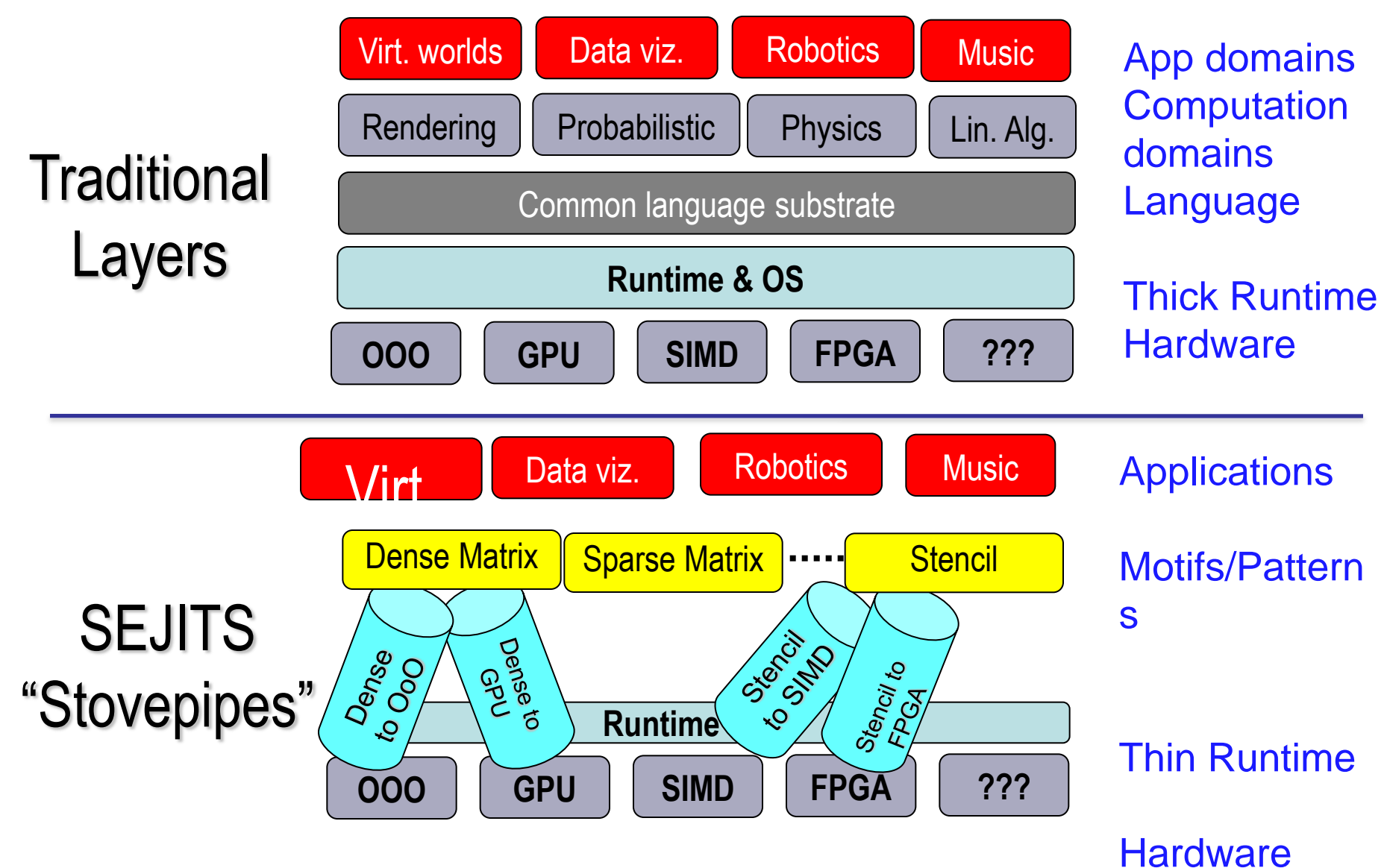
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SEJITS in a nutshell: Selective, Embedded Just-in-Time Specialization

- Productivity programmers write in *general purpose, modern, high level* PLL
- SEJITS infrastructure *specializes* computation patterns *selectively* at runtime
- Specialization uses runtime info to *generate* and *JIT-compile* ELL code targeted to hardware
- Embedded because PLL's own machinery enables (vs. extending PLL interpreter)

Layering vs. "Stovepiping"

- Layering: one or a few common intermediate languages
 - Must be flexible enough to support many DSLs
 - And map to wide variety of HW



- Stovepiping:** specialize structural computation patterns (*motifs*, not domains) directly to HW

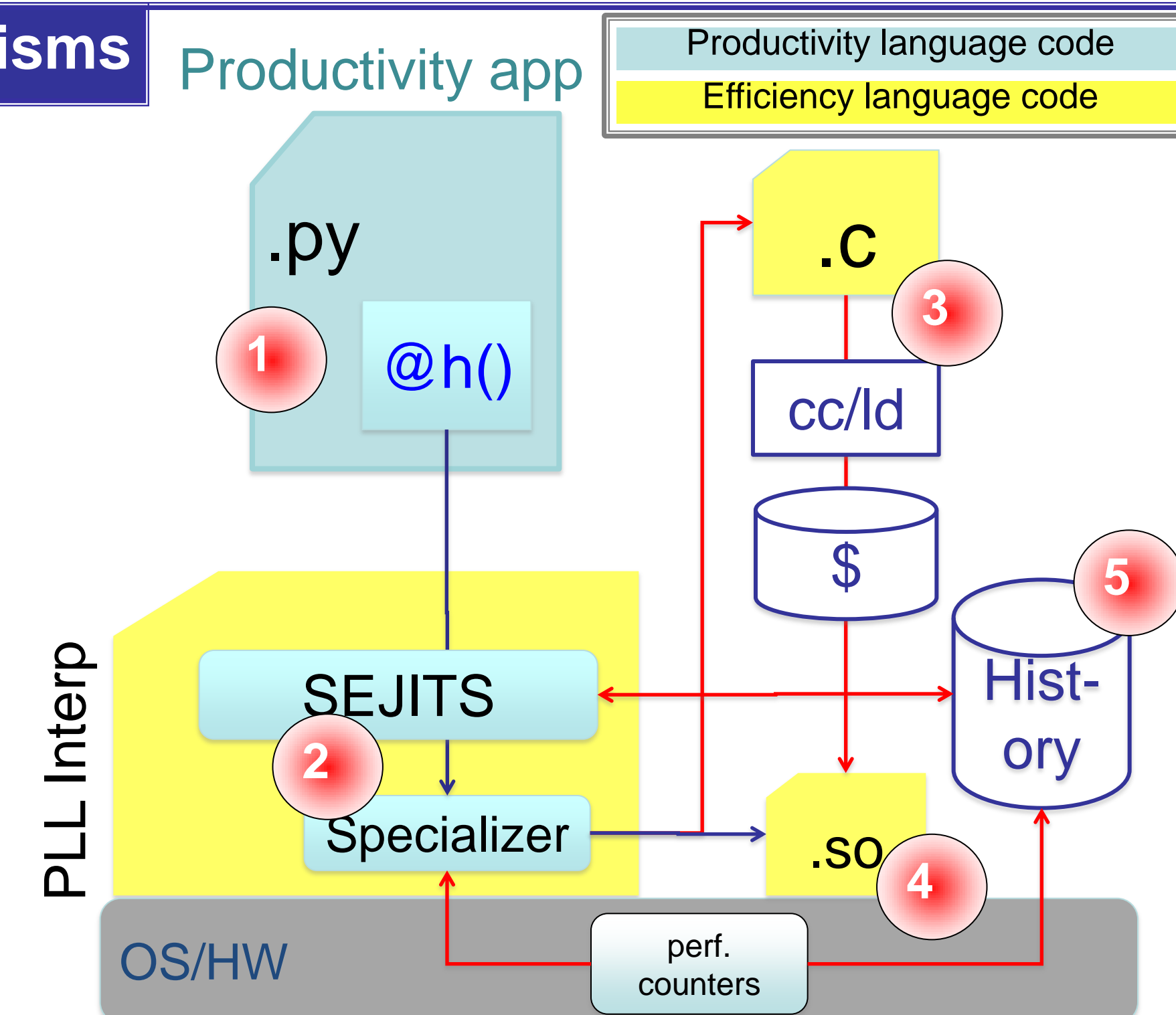
Leveraging Efficiency-Layer Research

- Efficiency programmers, autotuner writers: target *computation patterns* to hardware
 - stencil/SIMD codes => GPUs
 - sparse matrix => communication-avoiding algo's on multicore
 - "Big finance" Monte Carlo sim => MapReduce
- Libraries? Useful, but don't raise abstraction level
- How to make ELL work accessible to more PLL programmers?

SEJITS Exploits Productivity Level Language (PLL) Mechanisms

- Some functions in productivity app annotated as *potentially specializable*
- SEJITS intercepts calls using dynamic language features, uses introspection to examine function's Abstract Syntax Tree. If AST contains function call or pattern known in local catalog, *specializer* is invoked and handed AST
- Specializer generates source code in an efficiency language (C, OpenMP, CUDA, ...), compiles & links
- Specialized function binary is called, results returned to productivity language
- (Optional) performance recorded, code cached for future calls

Selective specialization: If any step fails, fall back to PLL (no need to JIT or specialize the whole app)



Embedded: SEJITS machinery uses PLL features, no need to modify or extend PLL interpreter

Early case study: Python + CUDA

- Ruby => OpenMP on multicore x86 (S. Kamil)
 - ~1000-2000x faster than pure Ruby
 - Minimal per-call overhead at runtime
- Python => NVidia GPU (B. Catanzaro, Y. Lee)
 - Stencils & Category-reduce (image processing)
 - Python decorators denote specializable functions
 - ~1000x Faster than pure Python
 - 3x-12x slower than handcrafted CUDA (including specialization overhead)
 - Overheads: Naive code generation & caching, Type propagation, CUDA compilation, data marshalling
- Productivity programmer only writes Python/Ruby, not CUDA or OpenMP

```
class LaplacianKernel < Kernel
def kernel(in_grid, out_grid)
  in_grid.each_interior do |point|
    in_grid.neighbors(point,1).each do |x|
      out_grid[point] += 0.2*x.val
    end
  end
end
```

```
VALUE kern_par(int argc, VALUE* argv, VALUE self) {
  unpack_arrays into in_grid and out_grid;
  private (t_6,t_7,t_8)
  for (t_8=1; t_8<256-1; t_8++) {
    for (t_7=1; t_7<256-1; t_7++) {
      for (t_6=1; t_6<256-1; t_6++) {
        int center = INDEX(t_6,t_7,t_8);
        out_grid[center] = (out_grid[center]
          +(0.2*in_grid[INDEX(t_6-1,t_7,t_8)]));
        ...
        out_grid[center] = (out_grid[center]
          +(0.2*in_grid[INDEX(t_6,t_7,t_8+1)]));
      }
    }
  }
  return Qtrue;}
}
```

Subverting PLL Mechanisms

- Observation: mechanisms intended to promote reuse also enable SEJITS
- Metaprogramming:** generate & JIT-compile efficiency code to replace PLL code for this function
 - Make decisions at runtime based on available HW, argument values, etc., vs. "static" autotuning
- Introspection:** intercept & analyze function to see if can specialize
 - Extend PLL without modifying interpreter
- Higher-order programming:** patterns at higher levels of abstraction
 - capture reusable *motifs* as well as low-level functions

Other Opportunities

- Autotuning
 - SEJITS can intercept calls and substitute autotuned code (see PySKI)
 - Locus of control for making co-tuning decisions
- Cloud Computing
 - Generate Hadoop (Java) code expressing PLL computation as MapReduce
 - Generate code for multiple cloud frameworks

Status, Ongoing Work, Challenges

- Prototypes working for NVidia, x86 multicore, RAMP (SPARC v8)
- Generalize infrastructure for catalog, pattern matching, call site annotation, history
- Integrate with PySKI/autotuning
- Cloud computing: Integrate with Nexus
- Cloud/multicore synergy: specialize intra-node as well as generate cloud code
- Capture additional motifs as specializers

Conclusions

- Enables code-generation strategy per-function, not per-app
- Uniform approach to productive programming
 - same app on cloud, multicore, autotuned libraries
- Research enabler
 - Incrementally develop specializers for different motifs, prototype HW
 - Don't need full compiler & toolchain just to get started