Performance Debugging: Methods and Tools

David Skinner

deskinner@lbl.gov

U.S. DEPARTMENT OF ENERGY | Office of Science

NERSC National Energy Research Scientific Computing Center

BERKELEY LAB Lawrence Berkeley National Laboratory

Performance Debugging: Methods and Tools

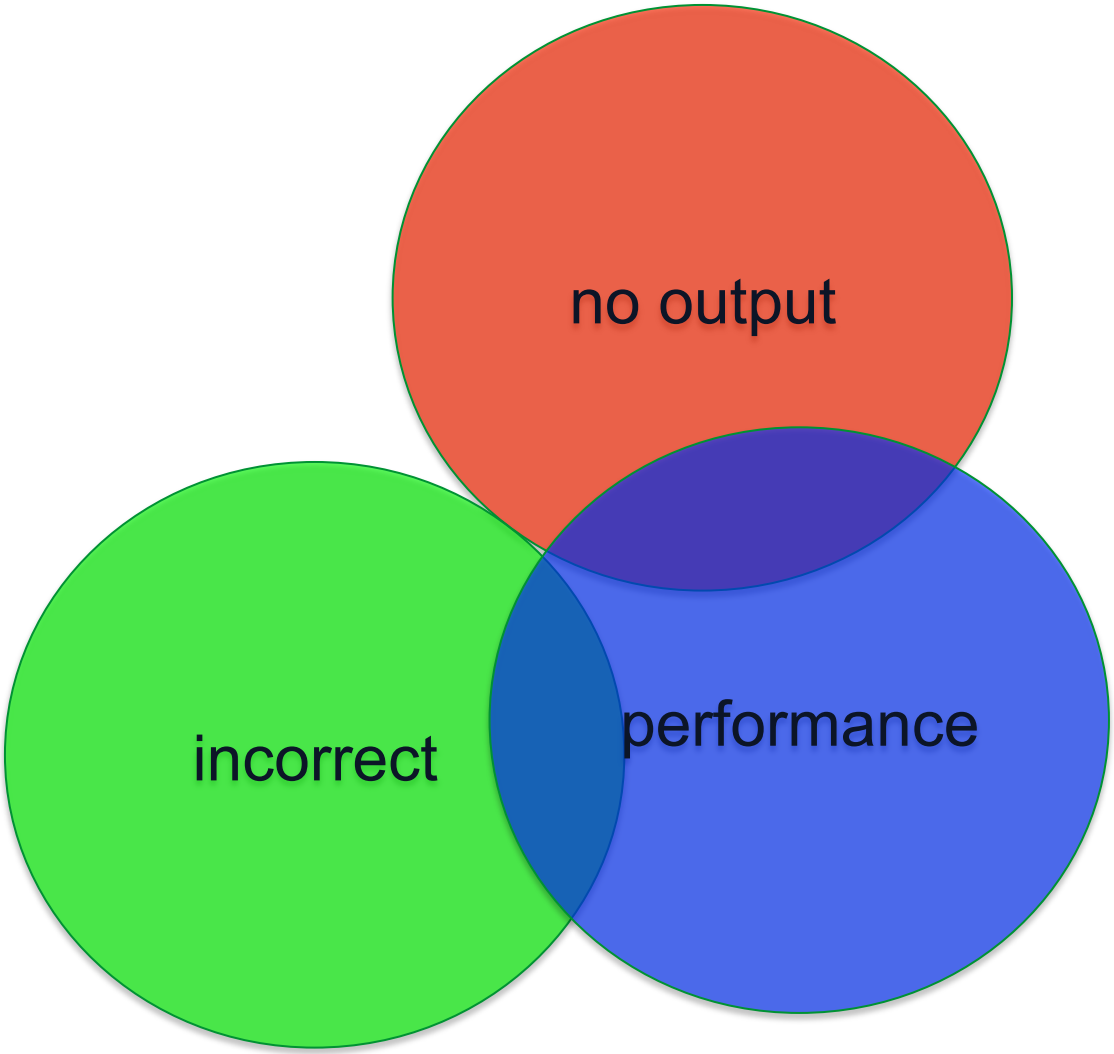
- **Principles**
  - Topics in performance scalability
  - Examples of areas where tools can help
- **Practice**
  - Where to find tools
  - Specifics to NERSC and Hopper
- **Scope & Audience**
  - Budding simulation scientist app dev
  - Compiler/middleware dev, YMMV

One Slide about NERSC

- **Serving all of DOE Office of Science**
  - domain breadth
- **Lots of users**
  - ~4K active
  - ~500 logged in

Big Picture of Performance and Scalability

What is performance?



Knight

July 2012

big fast computers

one misconfig

-$400M in 30min

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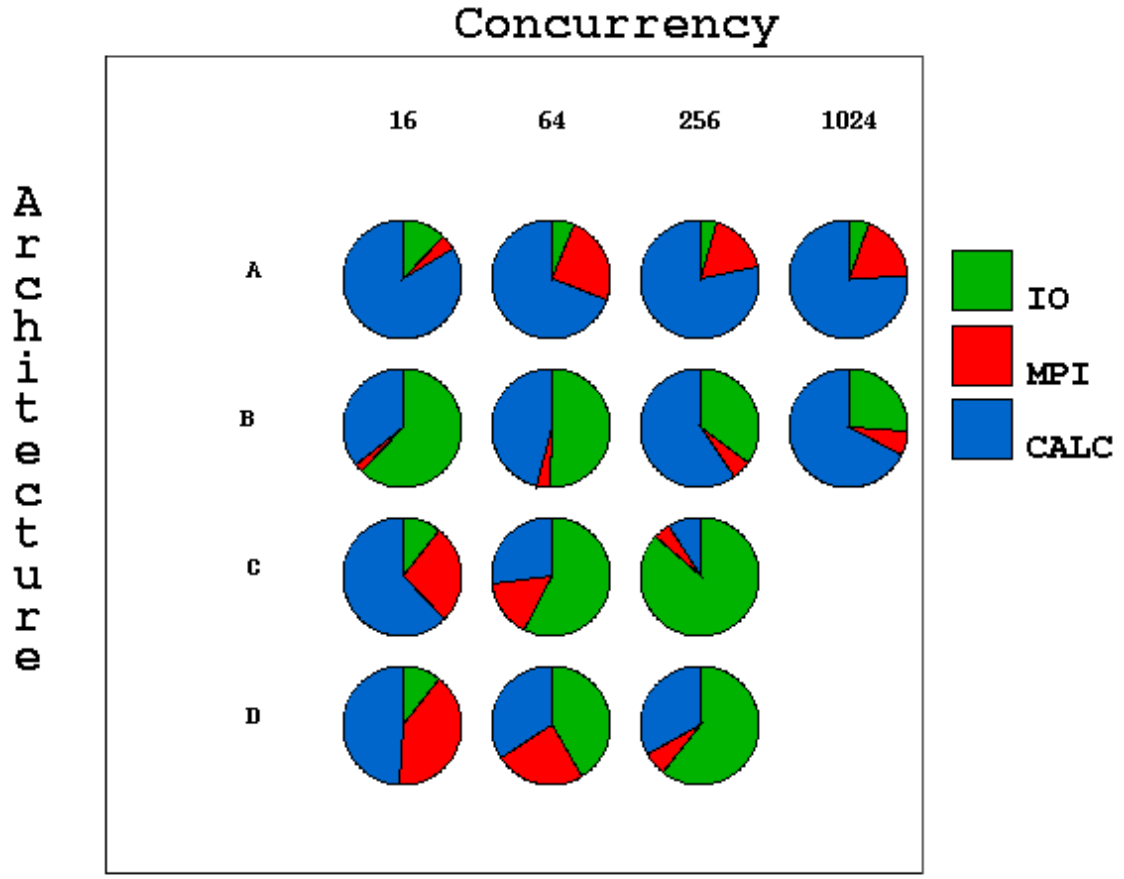
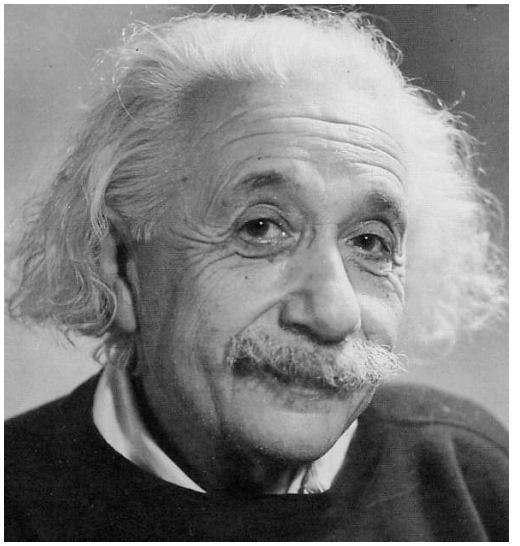
App Performance: Dimensions

- Code
- Input deck
- Computer
- Concurrency
- Workload
- Person



Performance is Relative

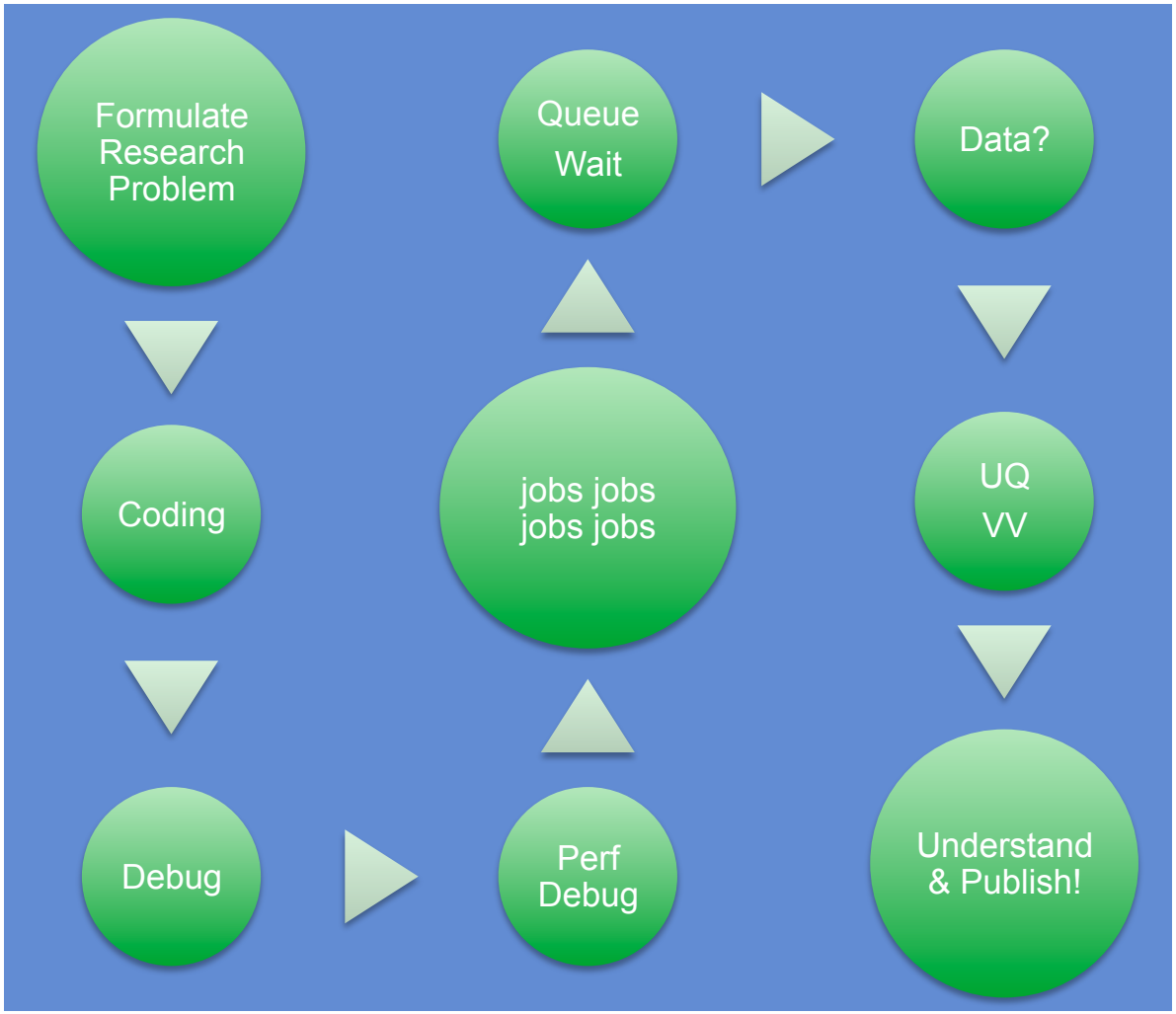
- **To your goals**
  - Time to solution, $T_q + T_{wall}$ …
  - Your research agenda
  - Efficient use of allocation



- **To the**
  - application code
  - input deck
  - machine type/state

Suggestion:
Focus on specific use cases
as opposed to making
everything
perform well.
Bottlenecks can shift.

Performance, more than a single number

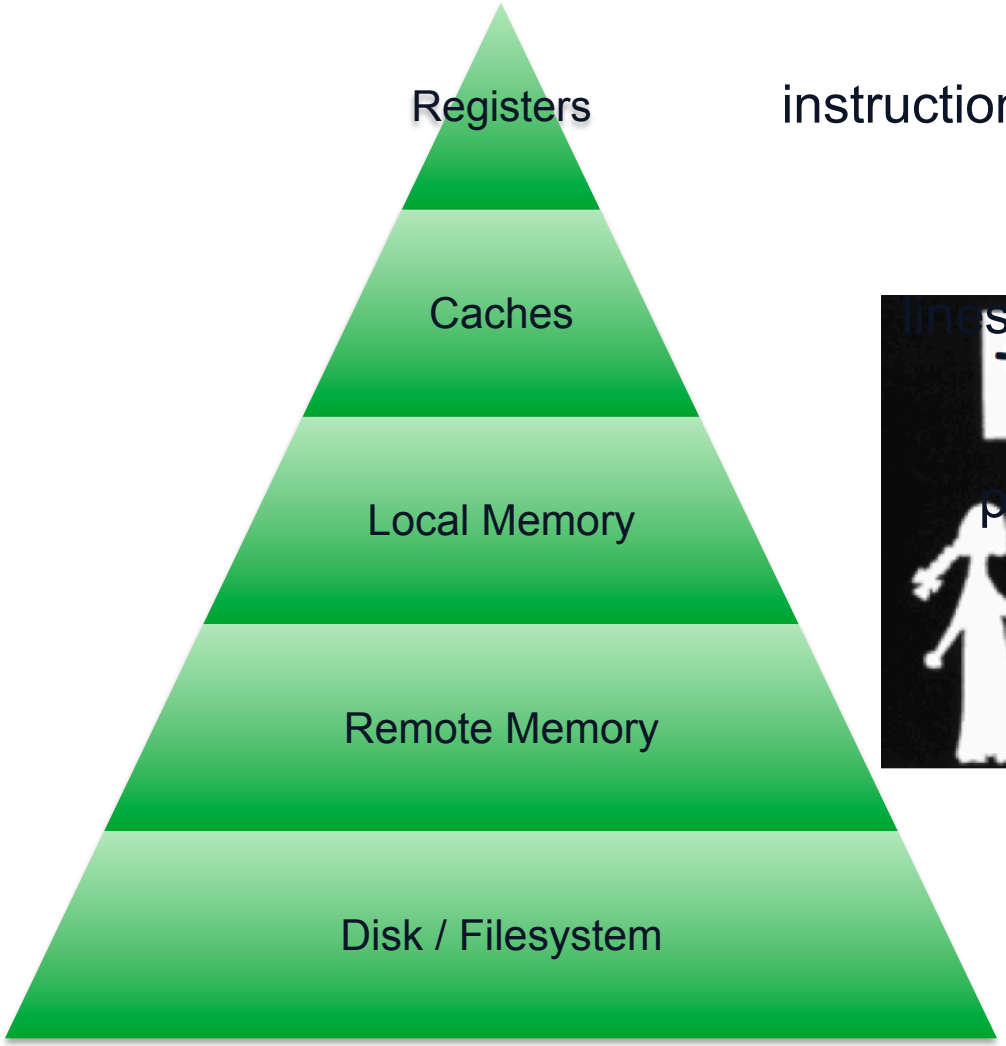


- Plan where to put effort
- Optimization in one area can de-optimize another
- Timings come from timers and also from your calendar, time spent coding
- Sometimes a slower algorithm is simpler to verify correctness

Different Facets of Performance

- **Serial**
  - Leverage ILP on the processor
  - Feed the pipelines
  - Exploit data locality
  - Reuse data in cache
- **Parallel**
  - Expose concurrency
  - Minimizing latency effects
  - Maximizing work vs. communication

Performance is Hierarchical



instructions & operands

lines

pages

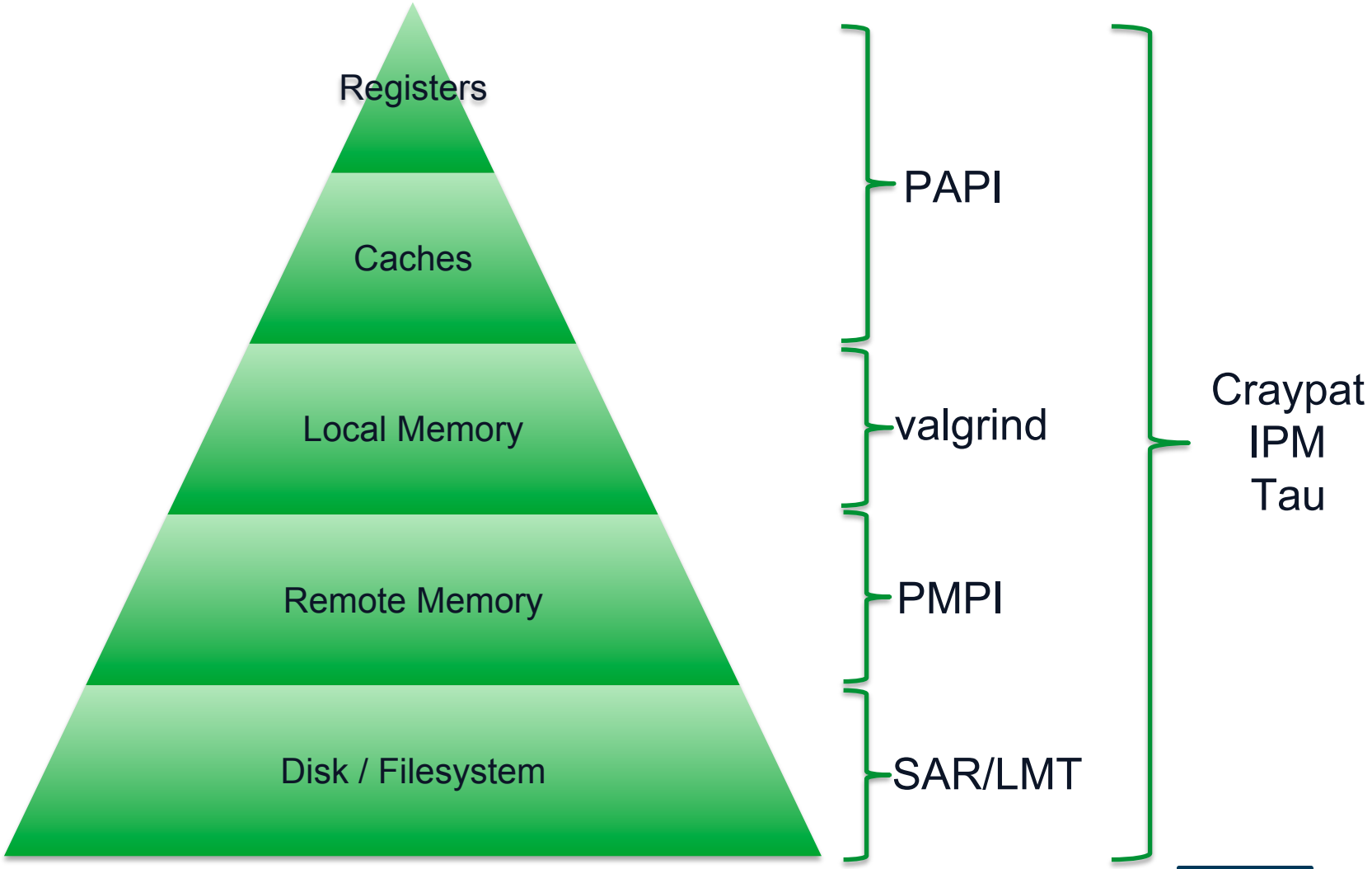
messages

blocks, files

…on to specifics about HPC tools

Mostly at NERSC but fairly general

Tools are Hierarchical



HPC Perf Tool Mechanisms

- **Sampling**
  - Regularly interrupt the program and record where it is
  - Build up a statistical profile
- **Tracing / Instrumenting**
  - Insert hooks into program to record and time events, document everything
- **Use Hardware Event Counters**
  - Special registers count events on processor
  - E.g. floating point instructions
  - Many possible events
  - Only a few (~4 counters)

Typical Tool Use Requirements

- **(Sometimes) Modify your code with macros, API calls, timers**
- **Compile your code**
- **Transform your binary for profiling/ tracing with a tool**
- **Run the transformed binary**
  - A data file is produced
- **Interpret the results with a tool**

Performance Tools @ NERSC

- **Vendor Tools:**
  - CrayPat
- **Community Tools :**
  - TAU (U. Oregon via ACTS)
  - PAPI (Performance Application Programming Interface)
  - gprof
- **IPM: Integrated Performance Monitoring**

What HPC tools can tell us?

- **CPU and memory usage**
  - FLOP rate
  - Memory high water mark
- **OpenMP**
  - OMP overhead
  - OMP scalability (finding right # threads)
- **MPI**
  - % wall time in communication
  - Detecting load imbalance
  - Analyzing message sizes

Using the right tool

Tools can add overhead to code execution

- **What level can you tolerate?**

Tools can add overhead to scientists

- **What level can you tolerate?**

Scenarios:

- **Debugging a code that is “slow”**
- **Detailed performance debugging**
- **Performance monitoring in production**

Introduction to CrayPat

- **Suite of tools to provide a wide range of performance-related information**
- **Can be used for both sampling and tracing user codes**
  - with or without hardware or network performance counters
  - Built on PAPI
- **Supports Fortran, C, C++, UPC, MPI, Coarray Fortran, OpenMP, Pthreads, SHMEM**
- **Man pages**
  - intro_craypat(1), intro_app2(1), intro_papi(1)

Using CrayPat @ Hopper

1. **Access the tools**
   - `module load perftools`
2. **Build your application; keep .o files**
   - `make clean`
   - `make`
3. **Instrument application**
   - `pat_build ... a.out`
   - Result is a new file, `a.out+pat`
4. **Run instrumented application to get top time consuming routines**
   - `aprun ... a.out+pat`
   - Result is a new file XXXXX.xf (or a directory containing .xf files)
5. **Run pat_report on that new file; view results**
   - `pat_report XXXXX.xf > my_profile`
   - `vi my_profile`
   - Result is also a new file: XXXXX.ap2

Guidelines for Optimization

Derived metric	Optimization needed when*	PAT_RT_HWPC
Computational intensity	< 0.5 ops/ref	0, 1
L1 cache hit ratio	< 90%	0, 1, 2
L1 cache utilization (misses)	< 1 avg hit	0, 1, 2
L1+L2 cache hit ratio	< 92%	2
L1+L2 cache utilization (misses)	< 1 avg hit	2
TLB utilization	< 0.9 avg use	1
(FP Multiply / FP Ops) or (FP Add / FP Ops)	< 25%	5
Vectorization	< 1.5 for dp; 3 for sp	12 (13, 14)

* Suggested by Cray

Perf Debug and Production Tools

- **Integrated Performance Monitoring**
- **MPI profiling, hardware counter metrics, POSIX IO profiling**
- **IPM requires no code modification & no instrumented binary**
  - Only a “module load ipm” before running your program on systems that support dynamic libraries
  - Else link with the IPM library
- **IPM uses hooks already in the MPI library to intercept your MPI calls and wrap them with timers and counters**

IPM: Let's See

1) Do "module load ipm", link with $IPM, then run normally

2) Upon completion you get

```
##IPM2v0.xx###################################################
#
# command   : ./fish -n 10000
# start     : Tue Feb 10 11:05:21 2012   host      : nid06027
# stop      : Tue Feb 10 11:08:19 2012   wallclock : 177.71
# mpi_tasks : 25 on 2 nodes              %comm     : 1.62
# mem [GB]  : 0.24                       gflop/sec : 5.06
…
```

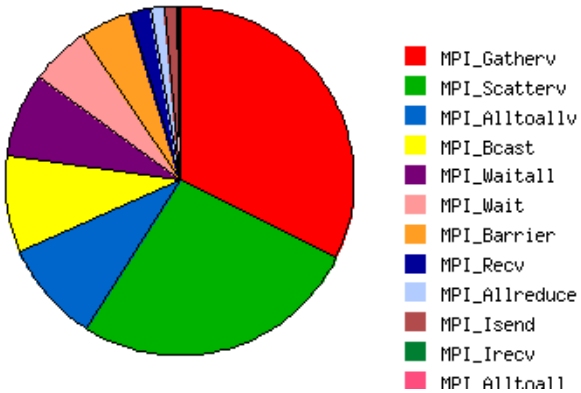
Maybe that's enough. If so you're done.

Have a nice day ☺

IPM : IPM_PROFILE=full

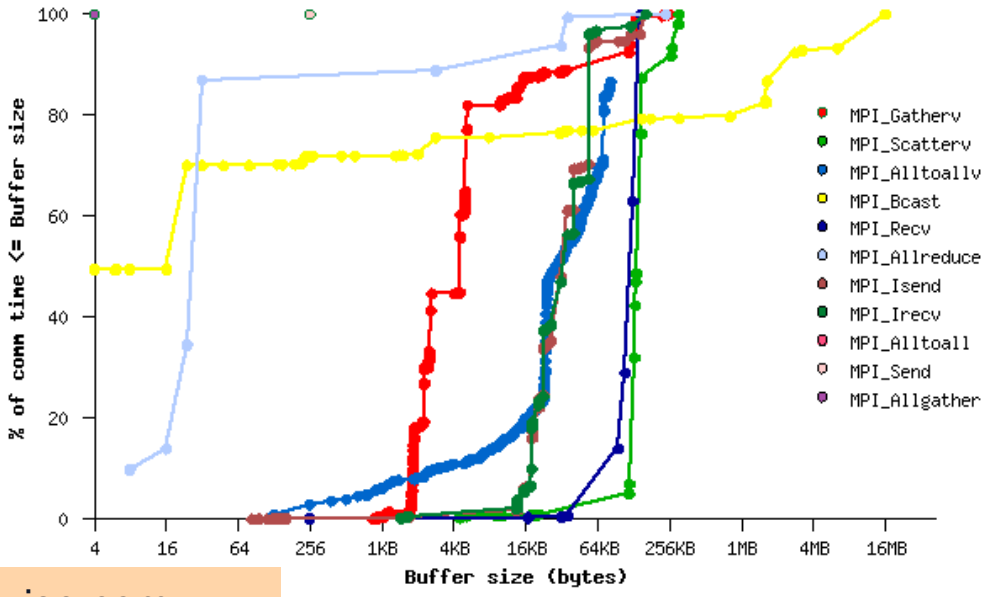
```
# host    : s05601/006035314C00_AIX        mpi_tasks : 32 on 2 nodes
# start   : 11/30/04/14:35:34              wallclock : 29.975184 sec
# stop    : 11/30/04/14:36:00              %comm     : 27.72
# gbytes  : 6.65863e-01 total              gflop/sec : 2.33478e+00 total
#                          [total]          <avg>           min           max
# wallclock                   953.272        29.7897       29.6092       29.9752
#user                         837.25         26.1641         25.71         26.92
# system                        60.6         1.89375          1.52          2.59
# mpi                         264.267         8.25834       7.73025       8.70985
# %comm                                       27.7234       25.8873       29.3705
# gflop/sec                   2.33478       0.0729619      0.072204     0.0745817
# gbytes                     0.665863       0.0208082     0.0195503     0.0237541
# PM_FPU0_CMPL           2.28827e+10    7.15084e+08   7.07373e+08   7.30171e+08
# PM_FPU1_CMPL           1.70657e+10    5.33304e+08   5.28487e+08   5.42882e+08
# PM_FPU_FMA             3.00371e+10     9.3866e+08   9.27762e+08   9.62547e+08
# PM_INST_CMPL           2.78819e+11    8.71309e+09   8.20981e+09   9.21761e+09
# PM_LD_CMPL             1.25478e+11    3.92118e+09   3.74541e+09   4.11658e+09
# PM_ST_CMPL             7.45961e+10    2.33113e+09   2.21164e+09   2.46327e+09
# PM_TLB_MISS            2.45894e+08    7.68418e+06   6.98733e+06   2.05724e+07
# PM_CYC                  3.0575e+11    9.55467e+09   9.36585e+09   9.62227e+09
#                            [time]       [calls]       <%mpi>      <%wall>
# MPI_Send                   188.386        639616         71.29         19.76
# MPI_Wait                   69.5032        639616         26.30          7.29
# MPI_Irecv                  6.34936        639616          2.40          0.67
# MPI_Barrier              0.0177442            32          0.01          0.00
# MPI_Reduce               0.00540609            32          0.00          0.00
# MPI_Comm_rank            0.00465156            32          0.00          0.00
# MPI_Comm_size           0.000145341            32          0.00          0.00
```

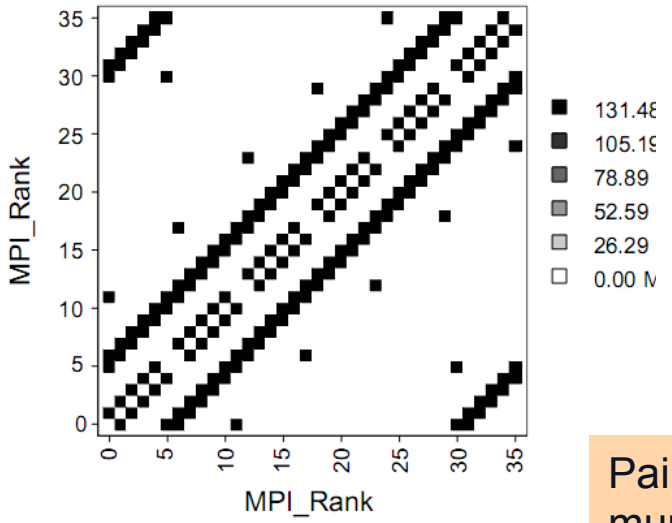
Analyzing IPM Data



Communication time per type of MPI call

CDF of time per MPI call over message sizes





Pairwise communication volume (comm. topology)

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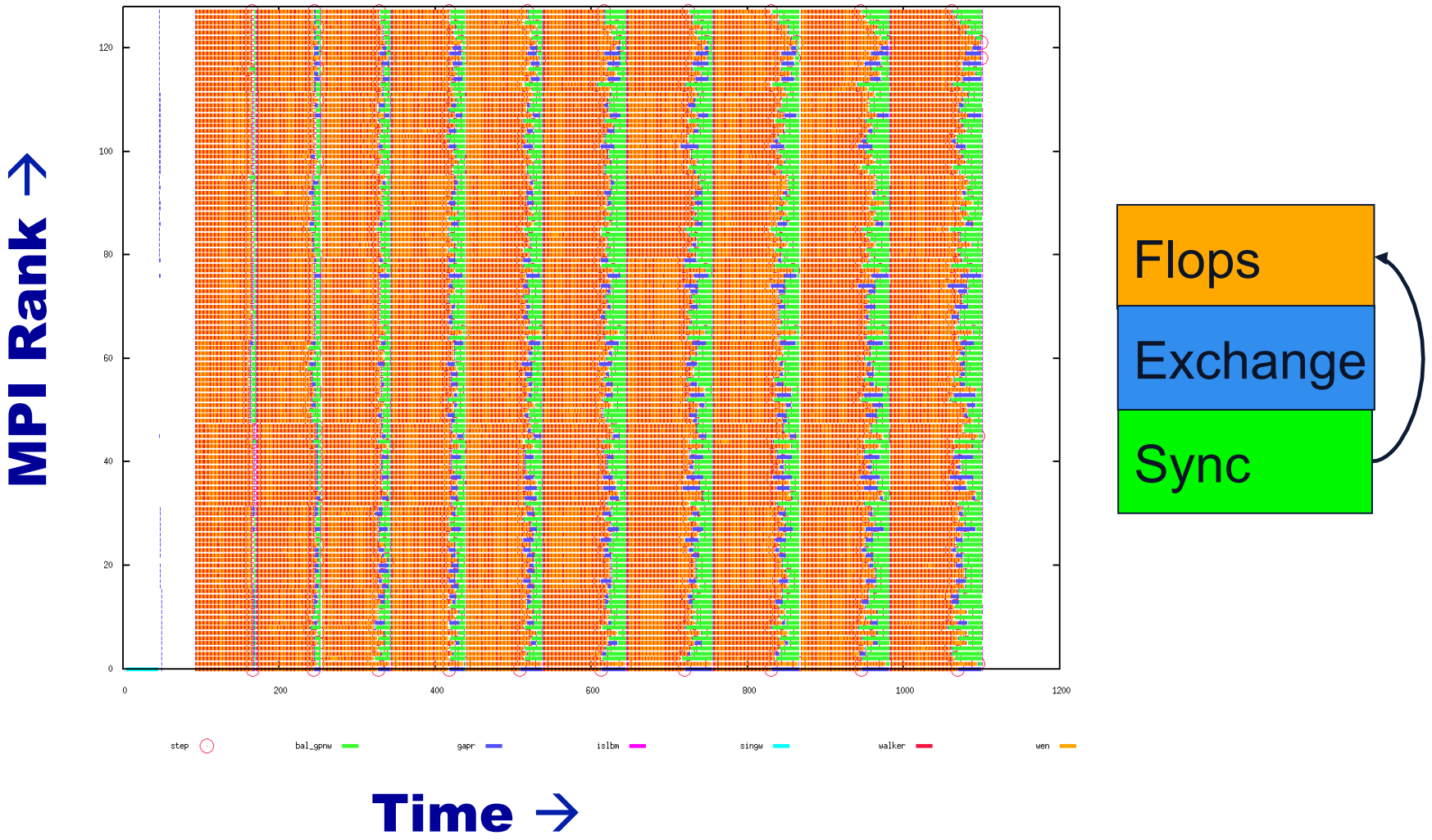
Tracing: Hard but Thorough



Advice: Develop (some) portable approaches to app optimization

- **There is a tradeoff between vendor-specific and vendor neutral tools**
  - Each have their roles, vendor tools can often dive deeper
- **Portable approaches allow apples-to-apples comparisons**
  - Events, counters, metrics may be incomparable across vendors
- **You can find printf most places**
  - Put a few timers in your code?

So you run your code with a perf tool and get some numbers…what do they mean?

Performance: Definitions

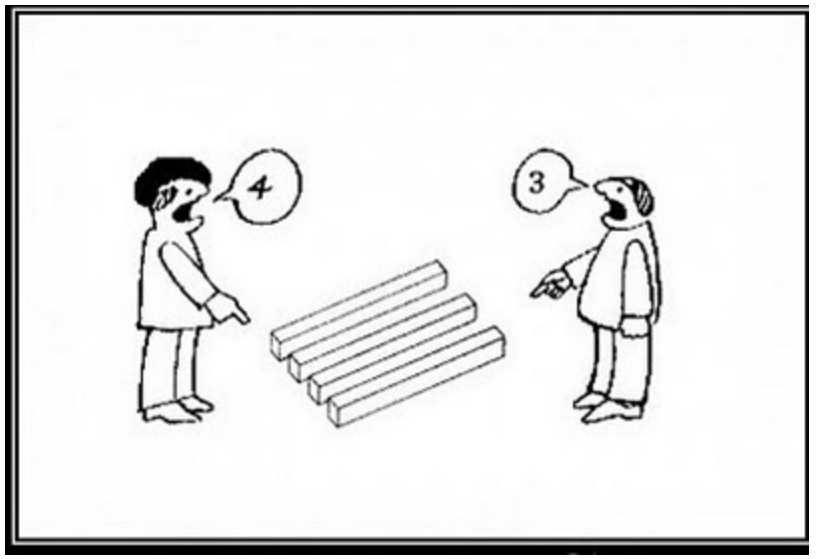
How do we measure performance?

Speed up = $T_s/T_p(n)$
Efficiency = $T_s/(n*T_p(n))$

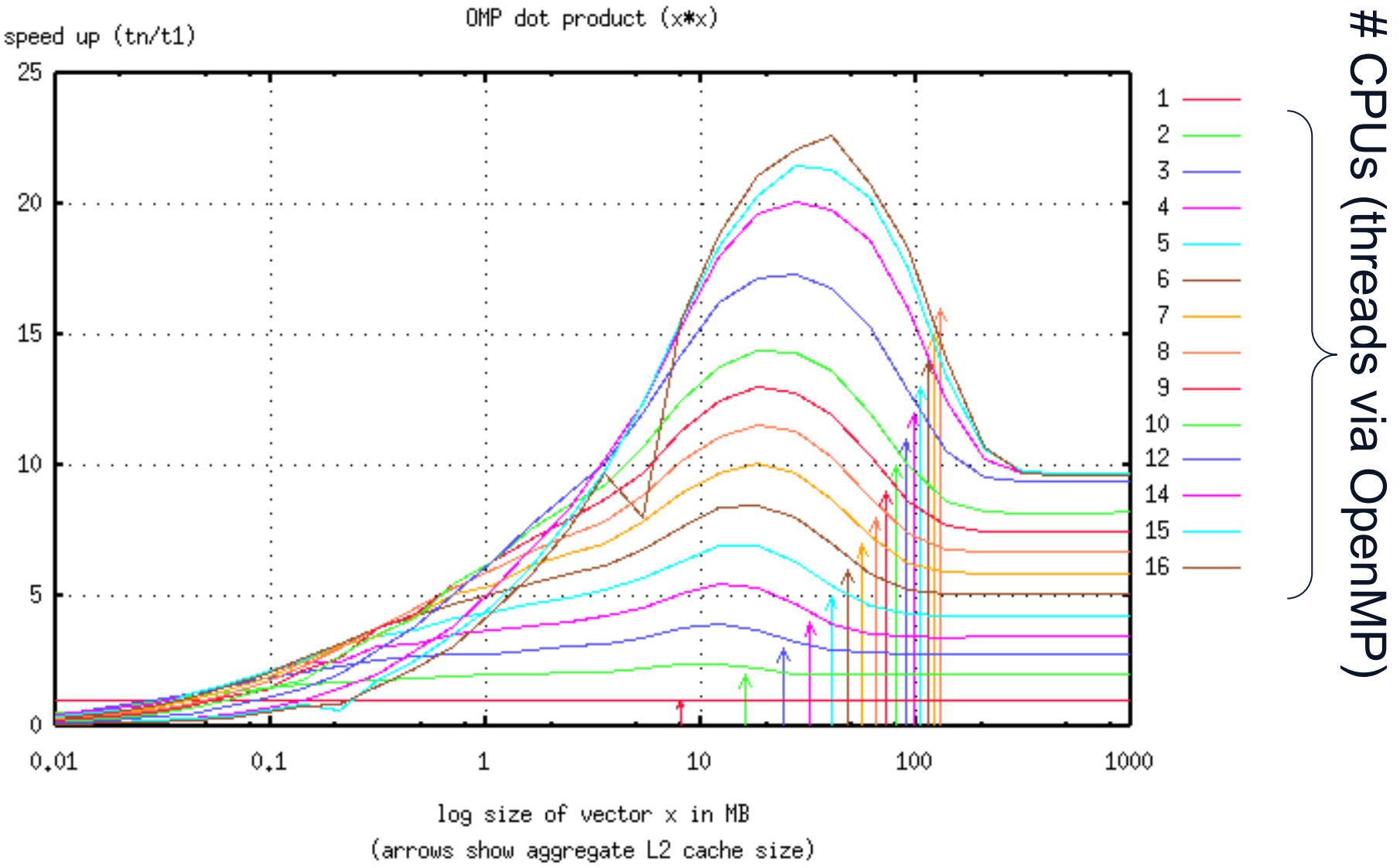
} Be aware there are multiple definitions for these terms

n=cores, s=serial, p=parallel

Isoefficiency: contours of constant efficiency amongst all problem sizes and concurrencies



Parallel speedups for x•x, how much?



Systematic Perf Measurement

- **Scaling studies involve changing the degree of parallelism. Will we be changing the problem also?**

- **Strong scaling :** Fixed problem size
- **Weak scaling:** Problem size grows with additional resources

- **Optimization: Is the concurrency chosing the problem size or vice versa?**

Sharks and Fish: Cartoon

Data:

n_fish is global

my_fish is local

$fish_i$ = {x, y, …}

Dynamics:

$$V \approx \sum_{ij} \frac{1}{r_{ij}}$$

$$F = ma$$

$$H = K + V$$

$$\dot{q} = -\partial H / \partial p$$

$$\dot{p} = -\partial H / \partial q$$

MPI_Allgatherv(myfish_buf, len[rank], .

```
for (i = 0; i < my_fish; ++i) {
      for (j = 0; j < n_fish; ++j) {  //  i!=j
       a_i += g * mass_j * ( fish_i – fish_j ) / r_ij
      }
}
```

Move fish

See a glimpse here:
http://www.leinweb.com/snackbar/wator/

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Sharks and Fish Results

Running on a NERSC machine

- **100 fish can move 1000 steps in**

  1 task → 5.459s
  32 tasks → 2.756s

  x 1.98 speedup

- **1000 fish can move 1000 steps in**

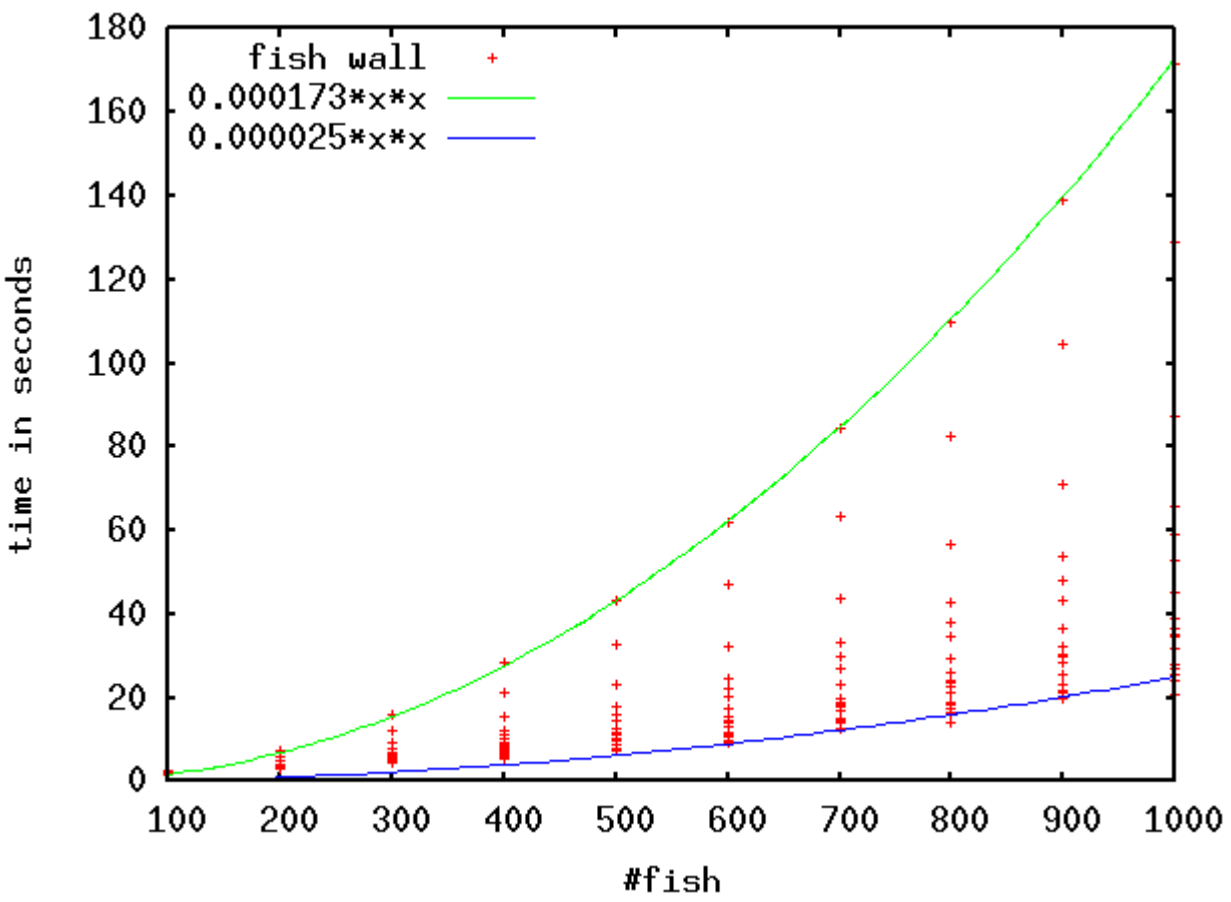
  1 task → 511.14s
  32 tasks → 20.815s

  x 24.6 speedup

- **What's the "best" way to run?**
  - How many fish do we really have?
  - How large a computer do we have?
  - How much "computer time" i.e. allocation do we have?
  - How quickly, in real wall time, do we need the answer?

Good 1st Step: Do runtimes make sense?

Running fish_sim for 100-1000 fish on 1-32 CPUs we see

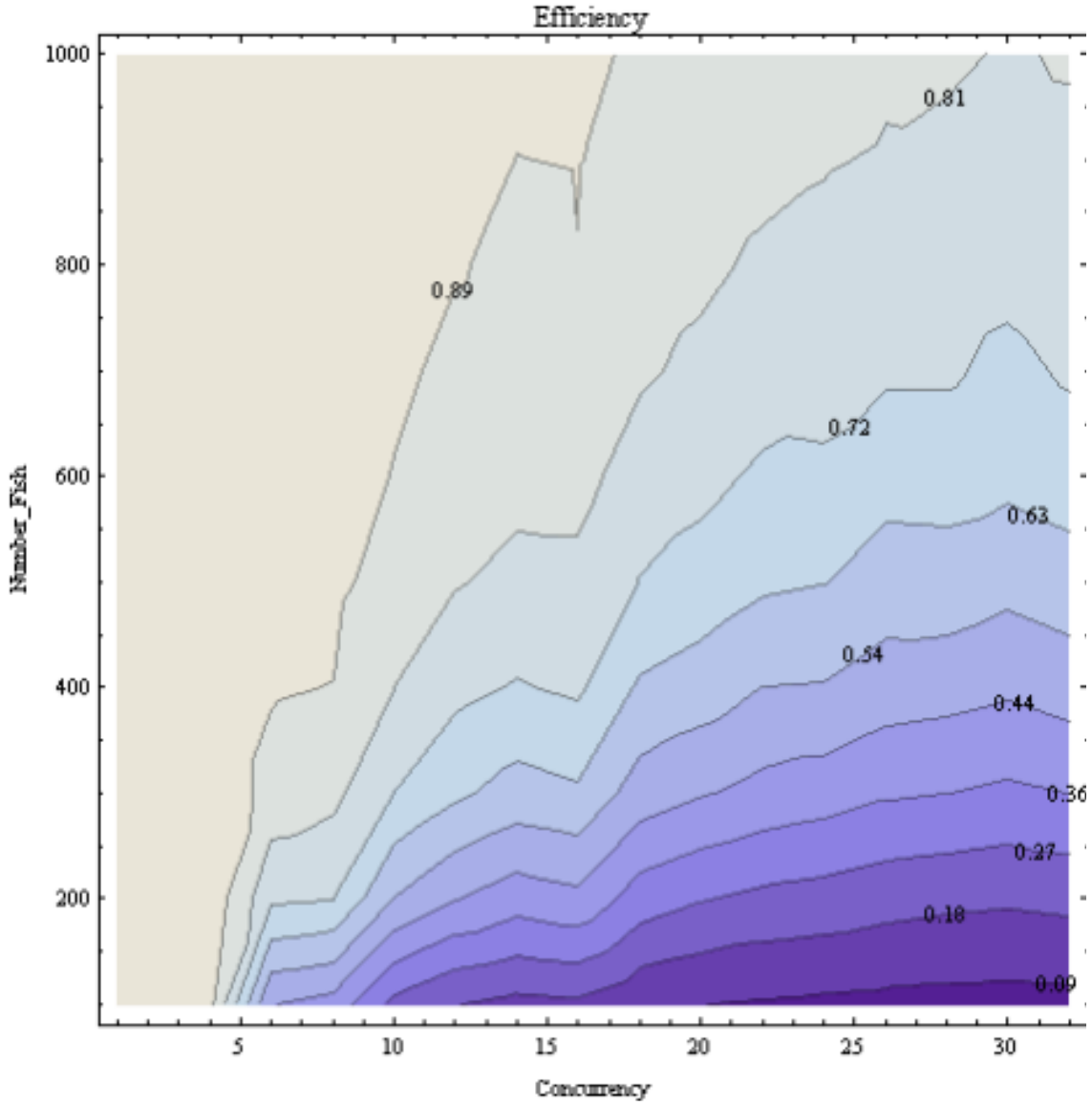


1 Task

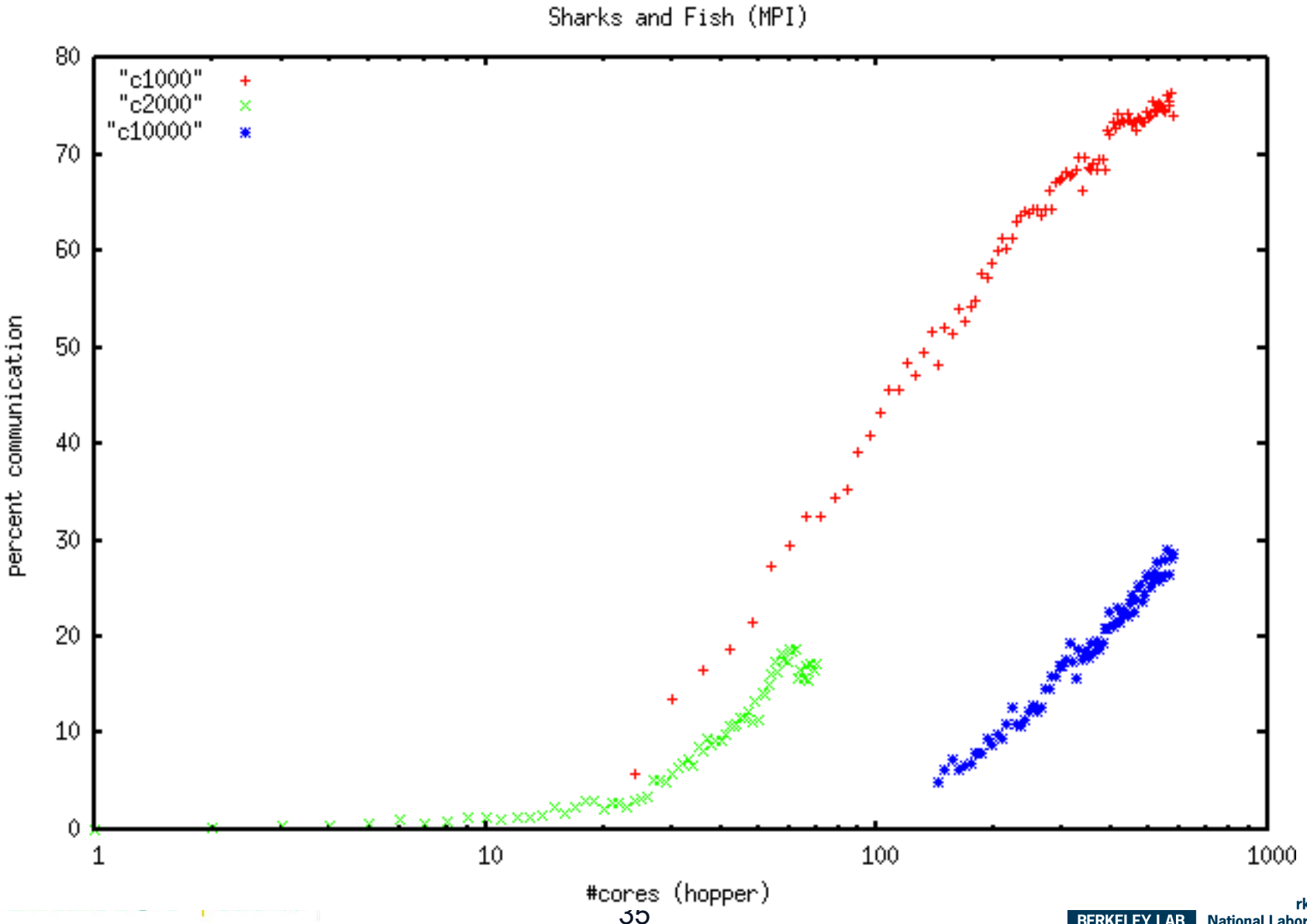
...

32 Tasks

Isoeffciencies



Too much communication



Scaling studies are not always so simple

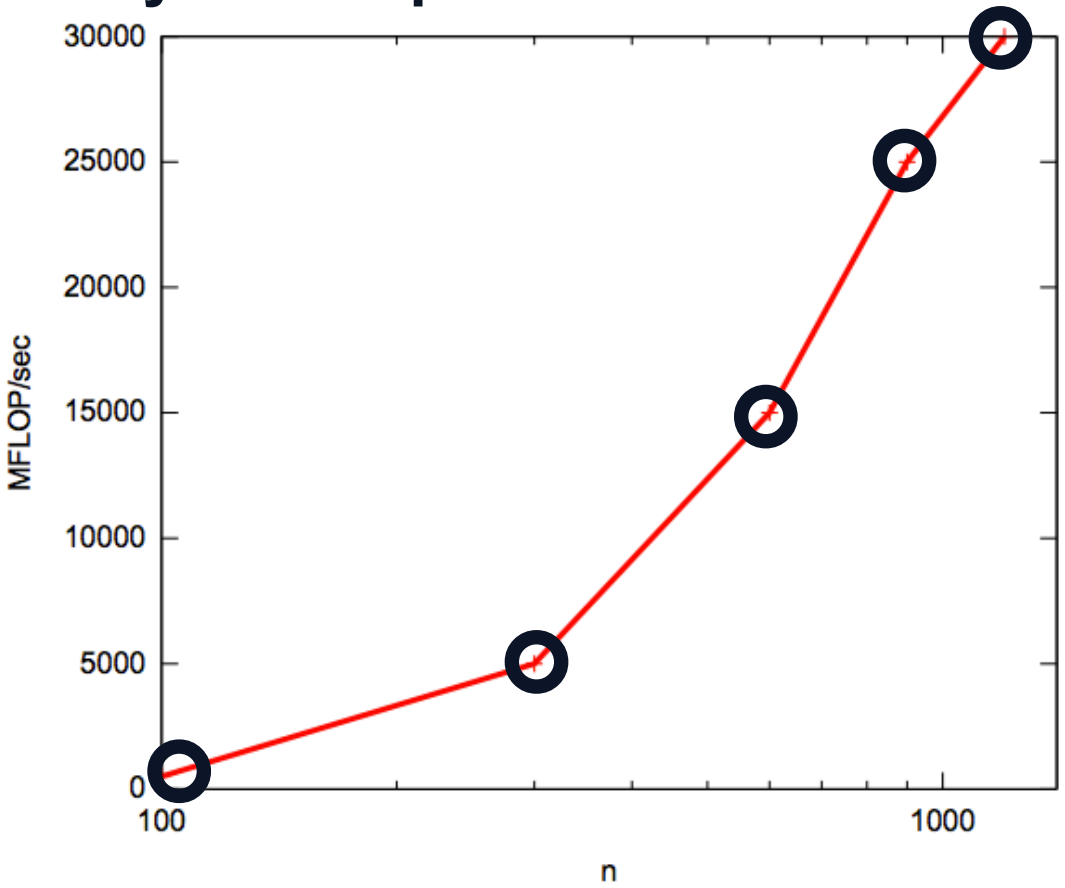
How many perf measurements?

With a particular goal in mind, we systematically vary concurrency and/or problem size

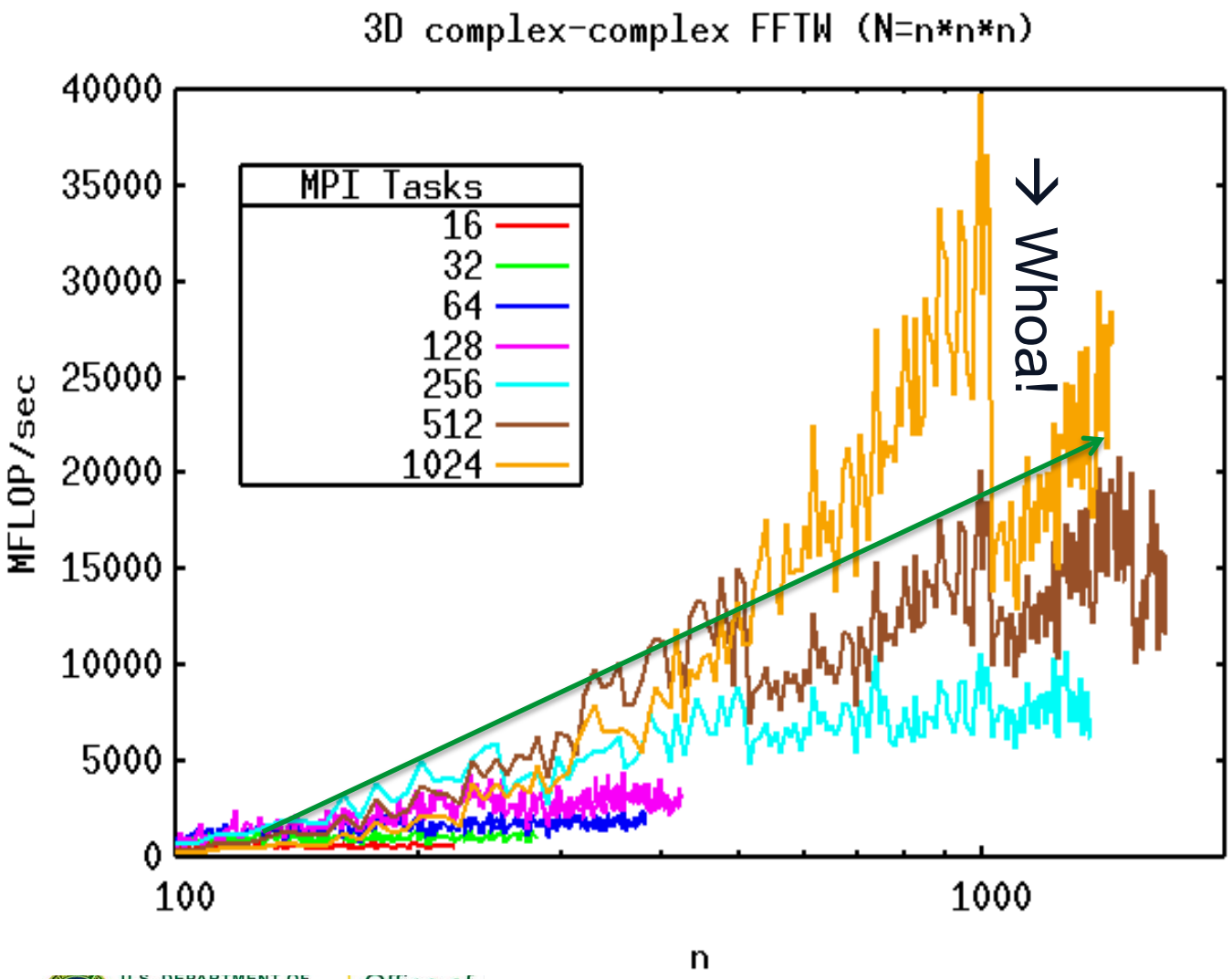
Example:

How large a 3D (n^3) FFT can I efficiently run on 1024 cpus?

Looks good?



The scalability landscape

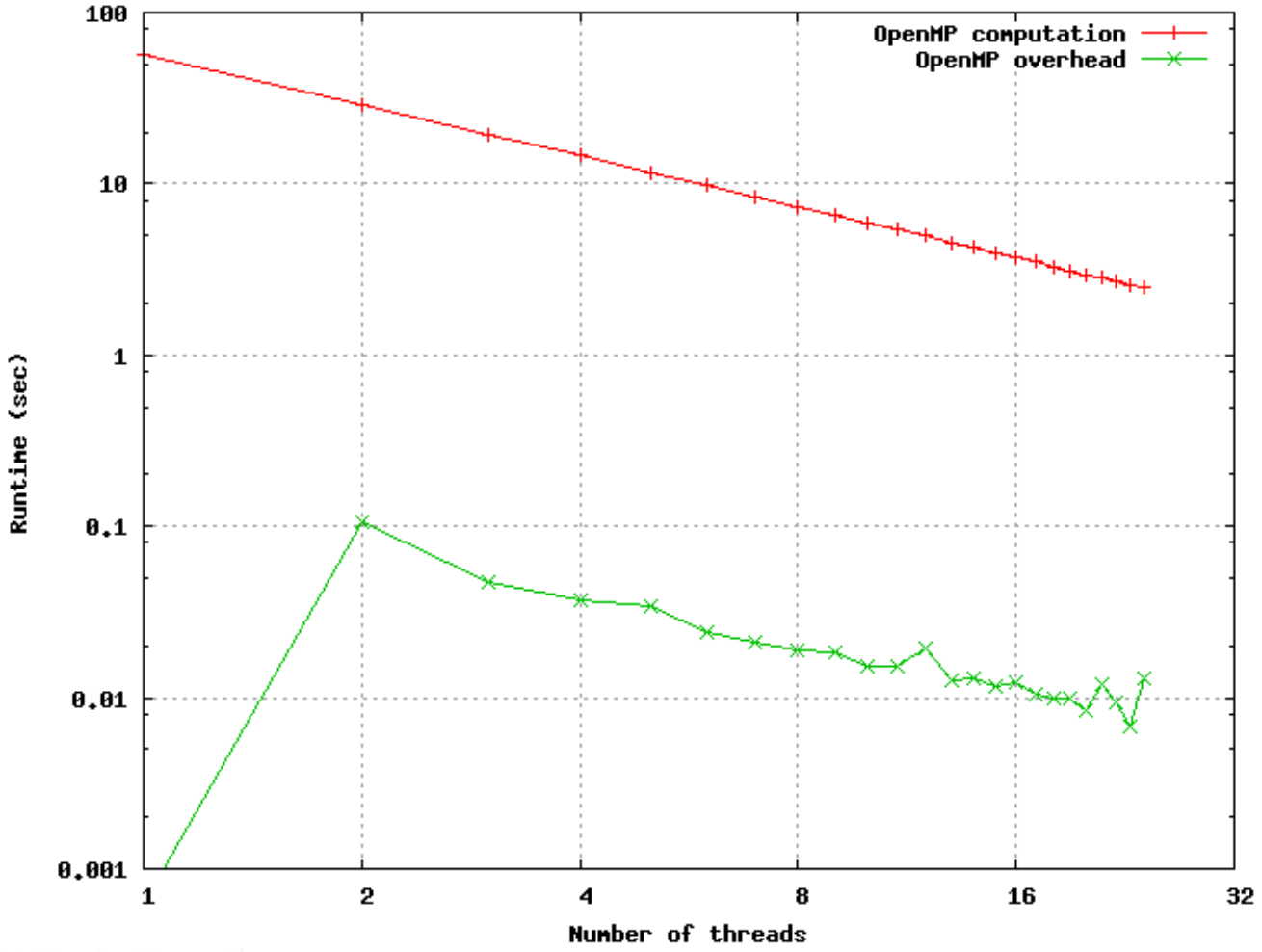


Why so bumpy?

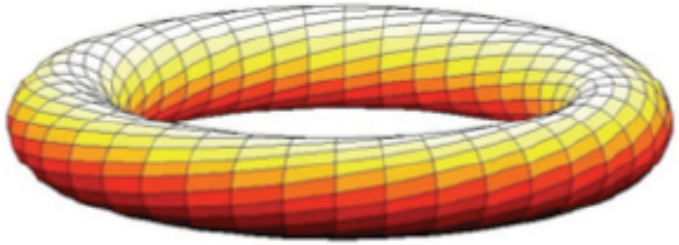
- Algorithm complexity or switching
- Communication protocol switching
- Inter-job contention
- ~bugs in vendor software

Not always so tricky

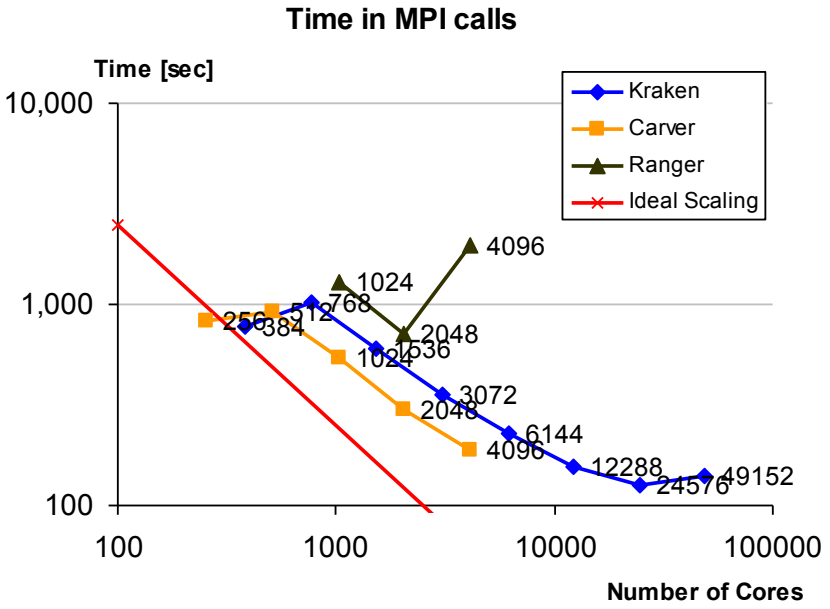
Main loop in jacobi_omp.f90; ngrid=6144 and maxiter=20

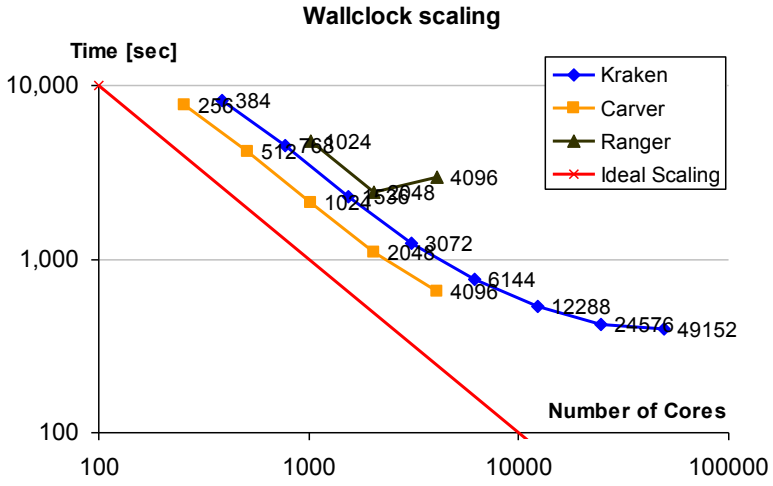


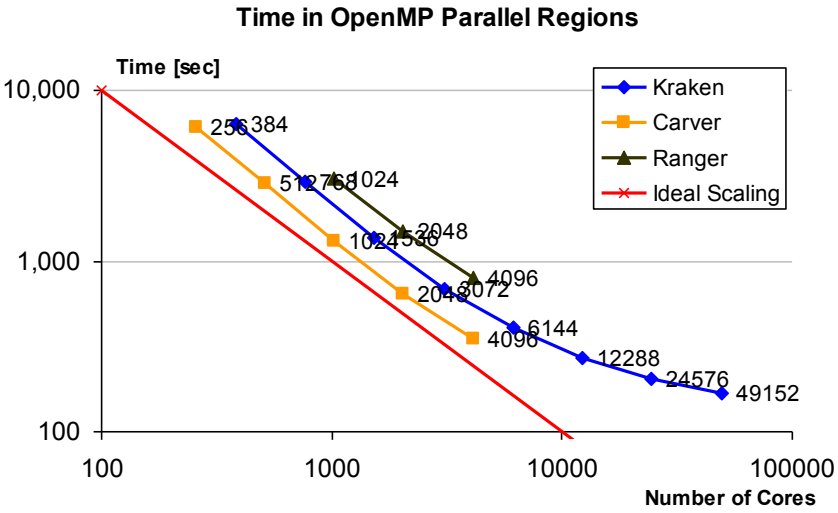
Scaling Studies at Scale



- Gyrokinetic Toroidal Code (fusion simulation)
  - OpenMP enabled 4/6 threads
  - Scaling up to 49152 cores
  - 3 machines







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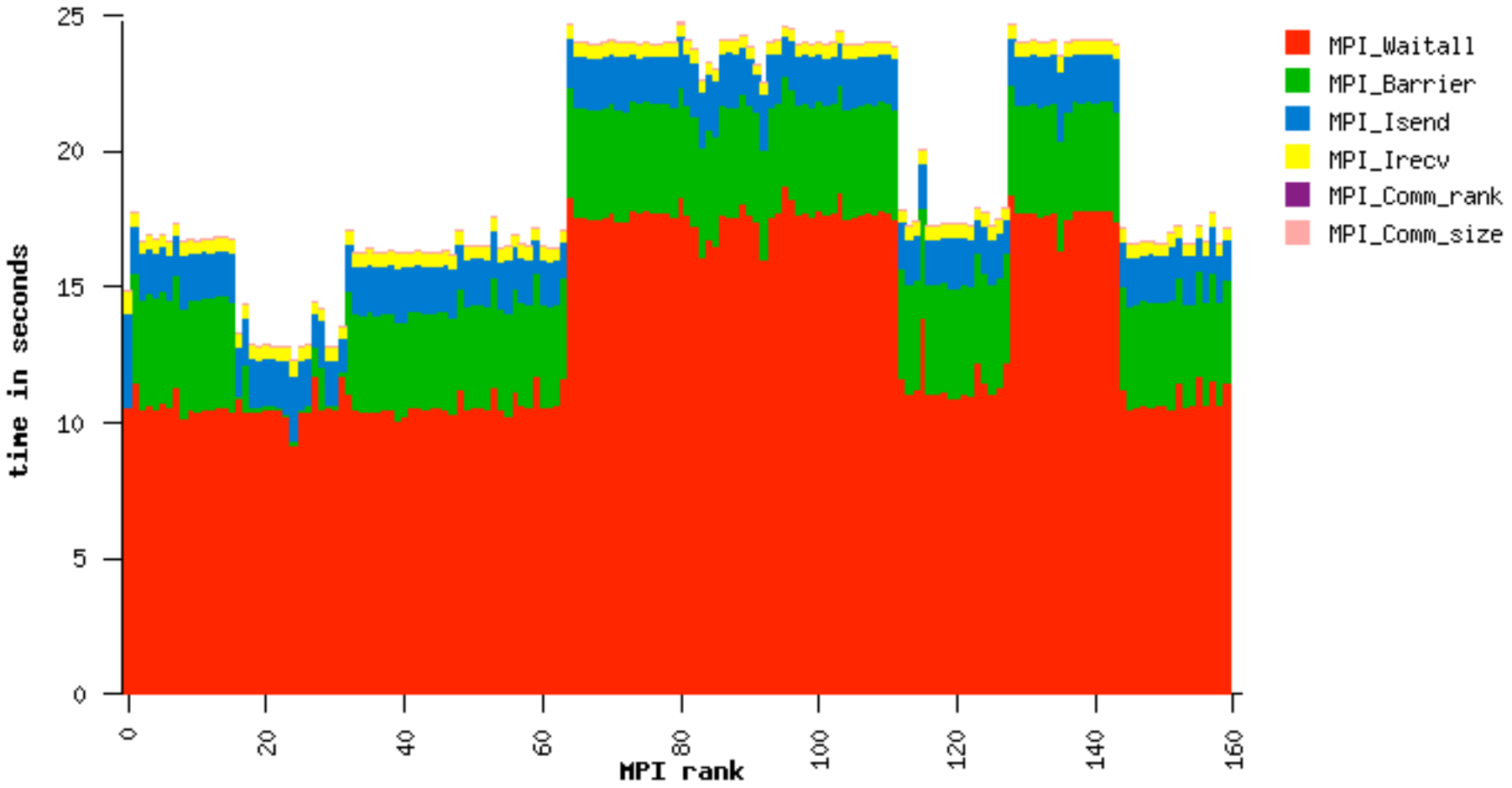
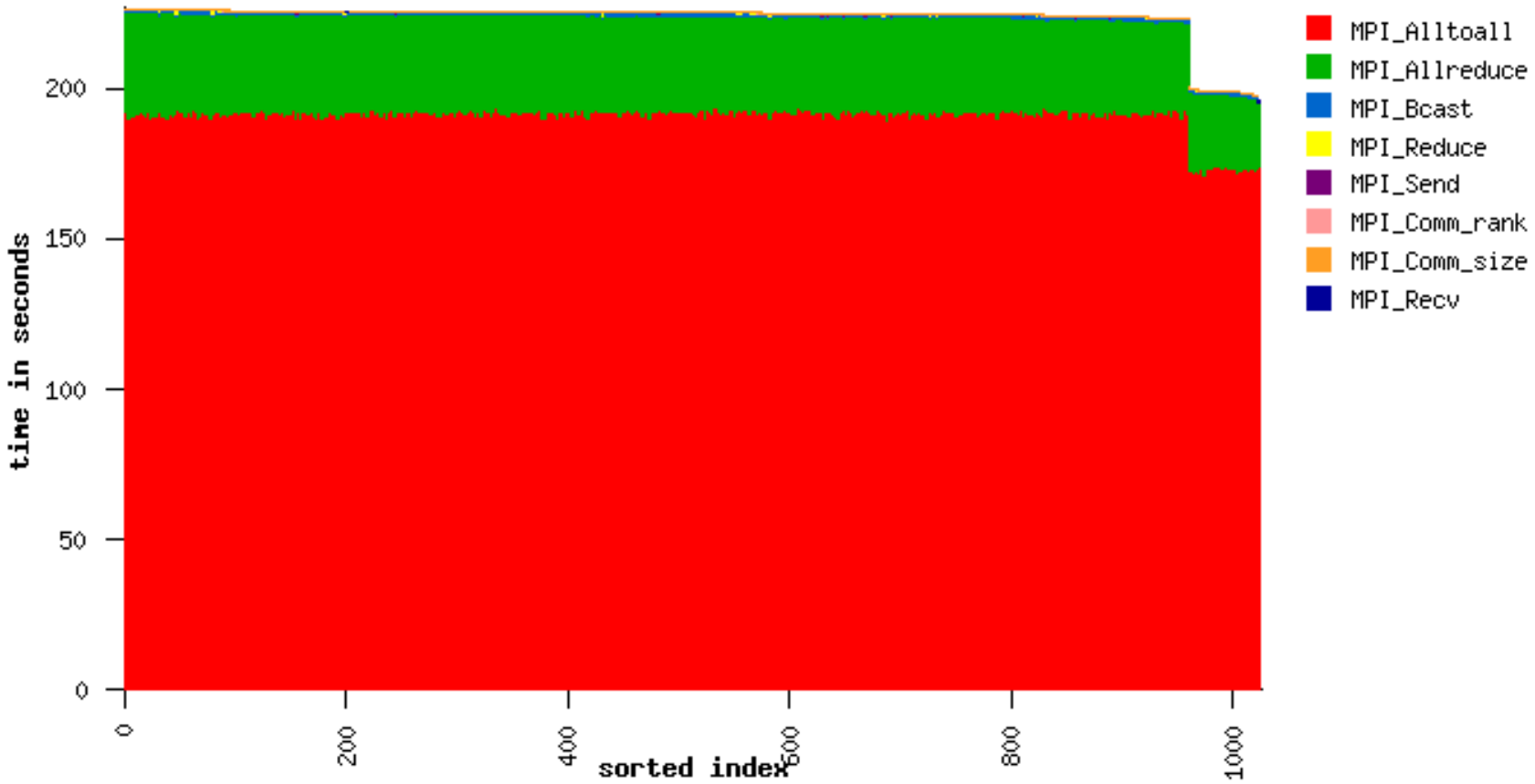
Let's look at scaling performance in more depth. A key impediment is often load imbalance



Load Imbalance : Pitfall 101

Communication Time: 64 tasks show 200s, 960 tasks show 230s



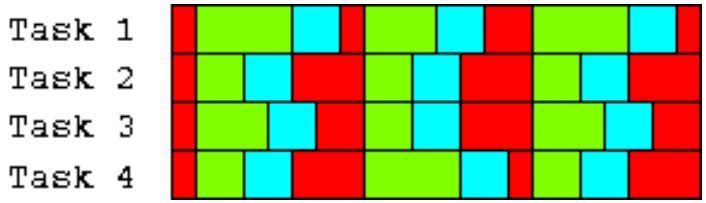
MPI ranks sorted by total communication time

Load imbalance is pernicious

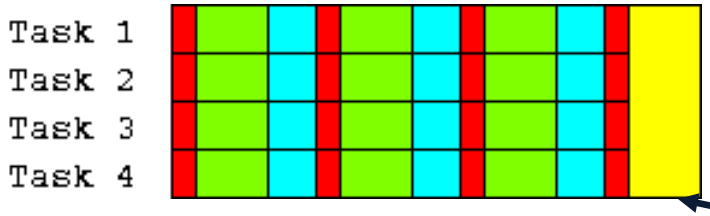
keep checking it

Load Balance : cartoon

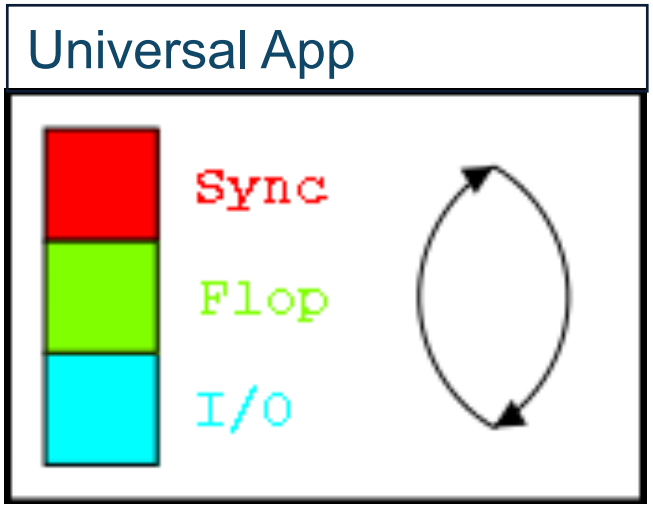
Unbalanced:



Balanced:

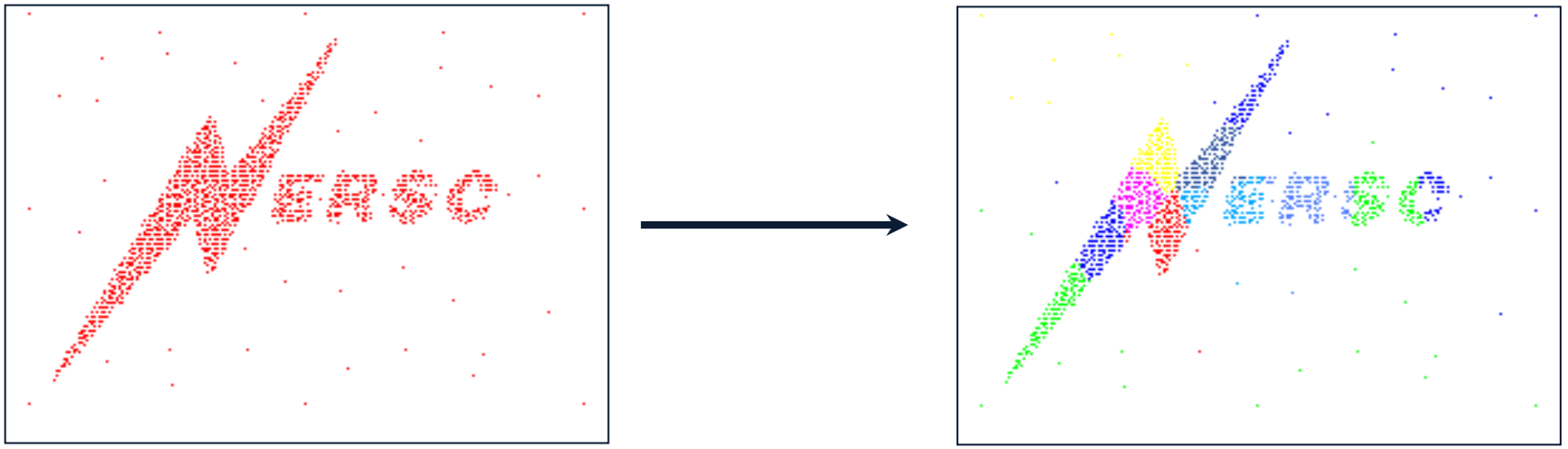


Time saved by load balance



Load Balance: Summary

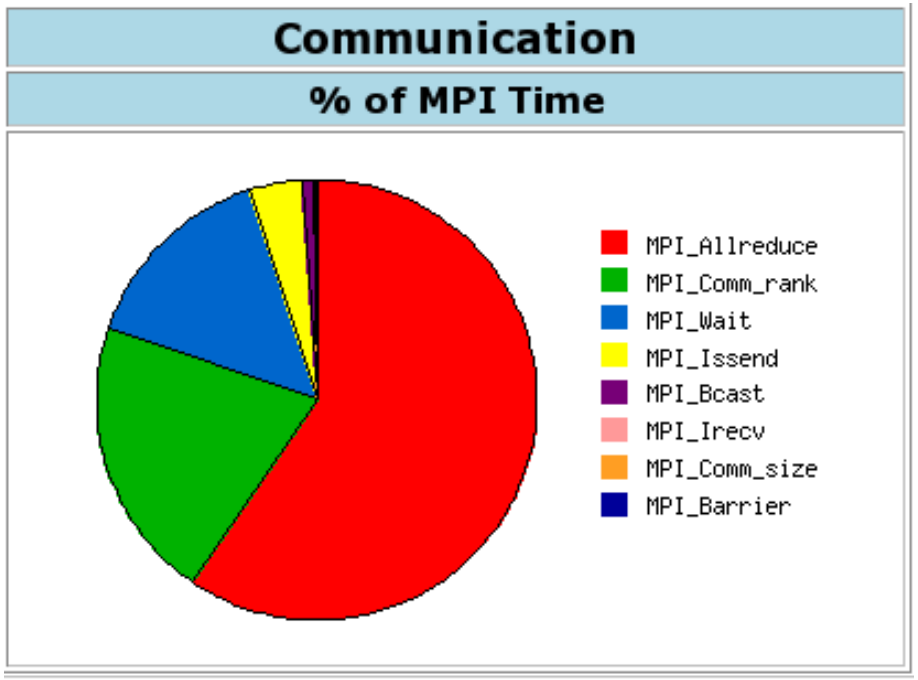
- Imbalance often a byproduct of
  1) data decomposition or 2) multi-core concurrency quirks
- Must be addressed **before** further MPI tuning can happen
- For regular grids consider padding or contracting
- Good software exists for graph partitioning / remeshing



- Dynamical load balance may be required for adaptive codes

Other performance scenarios

Simple Stuff: What's wrong here?



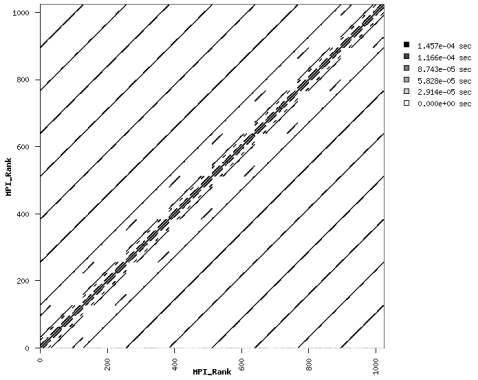
IPM Profile

This is why we need perf tools that are easy to use

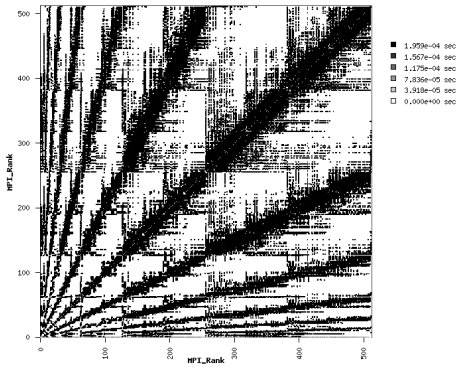
Communication Event Statistics (100.00% detail)

	Buffer Size	Ncalls	Total Time	Min Time	Max Time	%MPI	%Wall
MPI_Allreduce	8	3278848	124132.547	0.000	114.920	59.35	16.88
MPI_Comm_rank	0	35173439489	43439.102	0.000	41.961	20.77	5.91
MPI_Wait	98304	13221888	15710.953	0.000	3.586	7.51	2.14
MPI_Wait	196608	13221888	5331.236	0.000	5.716	2.55	0.72
MPI_Wait	589824	206848	5166.272	0.000	7.265	2.47	0.70

Not so simple: Comm. topology



MILC



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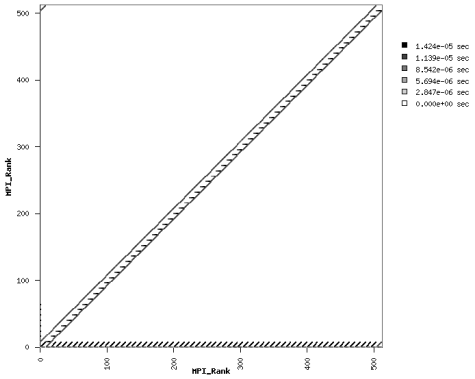

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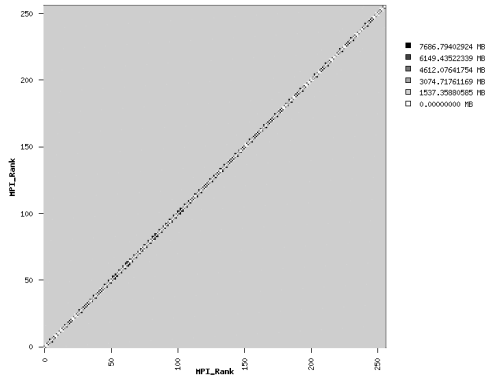

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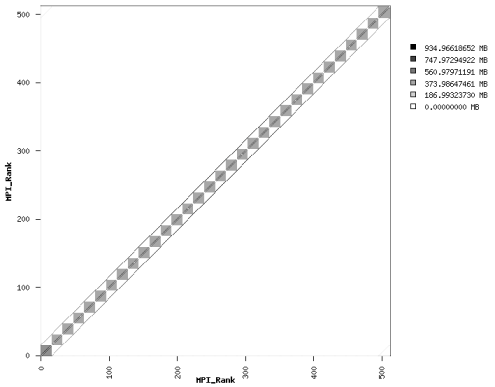

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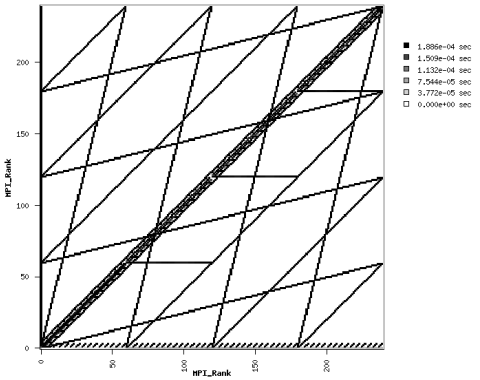

CAM

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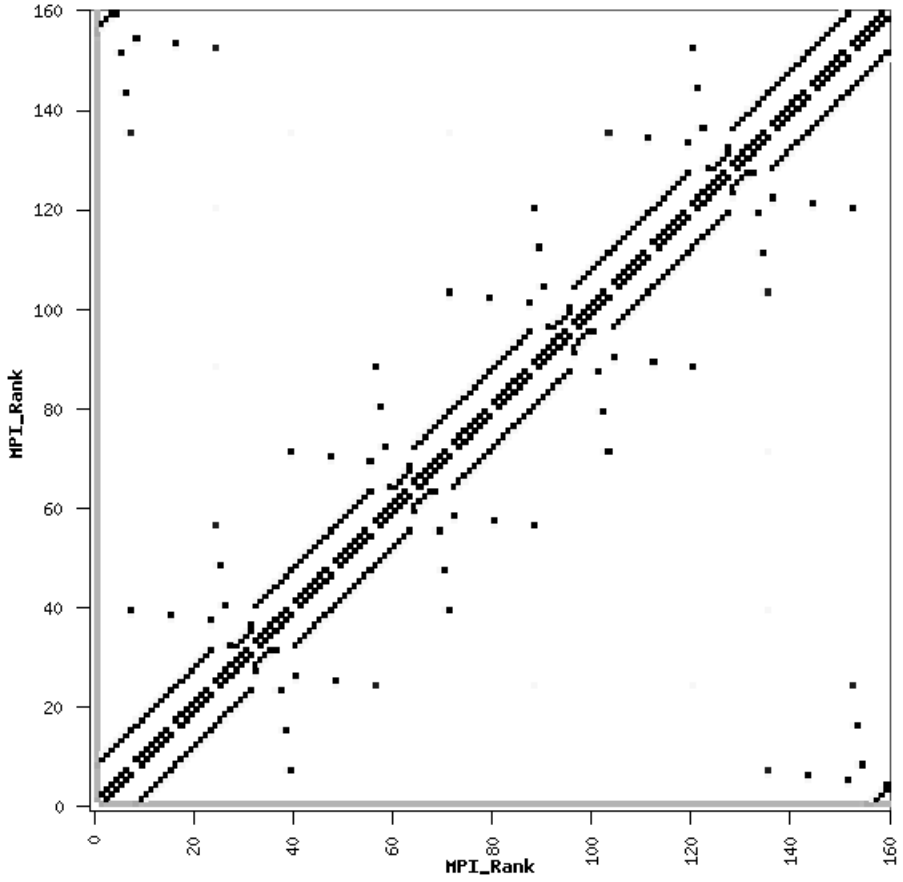
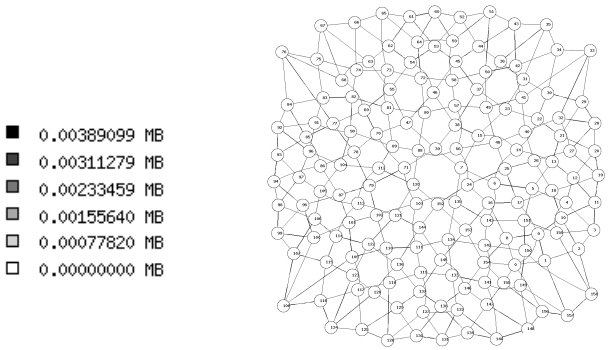
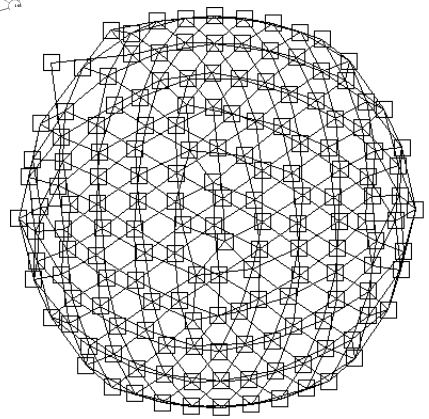
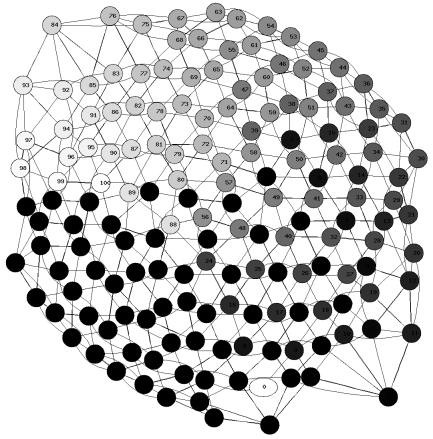
Application Topology









Performance in Batch Queue Space

A few notes on queue optimization

Consider your schedule

- **Charge factor**
  - regular vs. low
- **Scavenger queues**
- **Xfer queues**
  - Downshift concurrency

Consider the queue constraints

- **Run limit**
- **Queue limit**
- **Wall limit**
  - Soft (can you checkpoint?)

Jobs can submit other jobs

Marshalling your own workflow

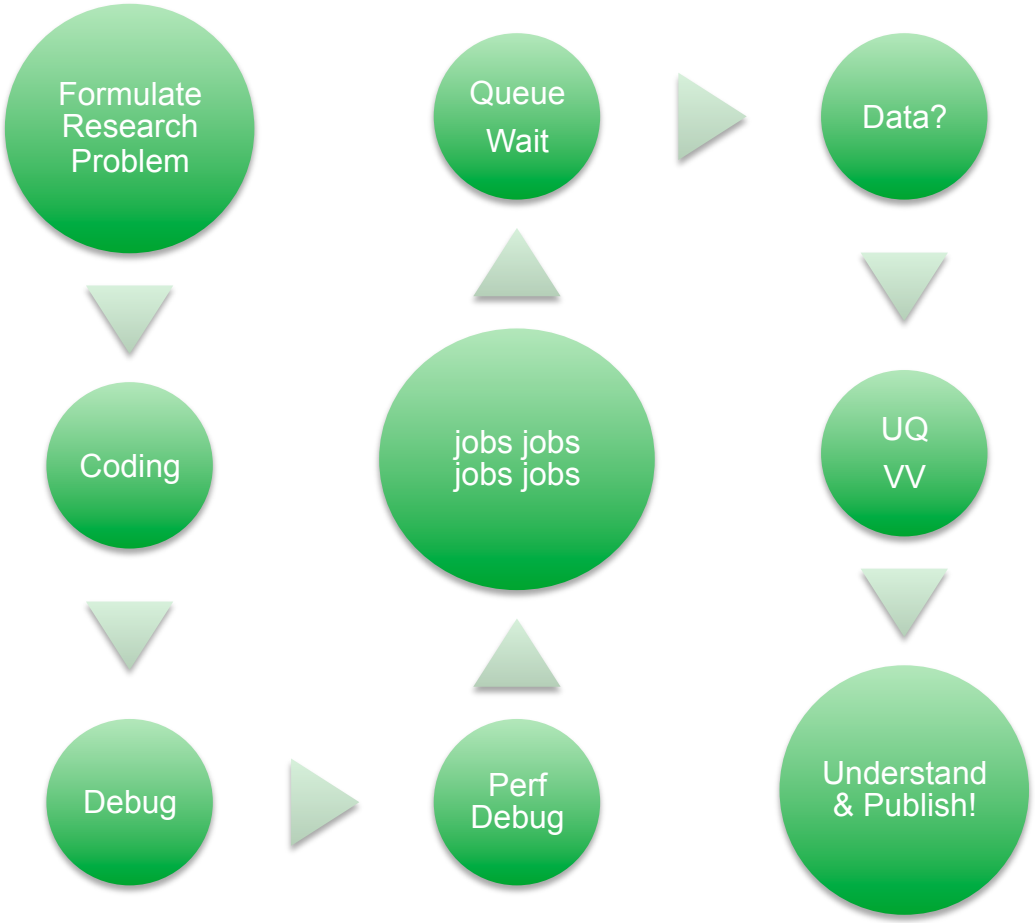
- **Lots of choices in general**
  - Hadoop, CondorG, MySGE
- **On hopper it's easy**

```
#PBS -l mppwidth=4096
aprun –n 512 ./cmd &
aprun –n 512 ./cmd &
…
aprun –n 512 ./cmd &

wait
```

```
#PBS -l mppwidth=4096
while(work_left) {
 if(nodes_avail) {
 aprun –n X next_job &
 }
wait
}
```

Thanks!



Contacts:

help@nersc.gov

deskinner@lbl.gov